



Exercise Session 01 – Asymptotics

Data Structures and Algorithms

These slides are based on those of the lecture, but were adapted and extended by the teaching assistant Adel Gavranović

Today's Schedule

Intro

Learning Objectives

Exercise Process

Repetition Theory

Examples (Theory)

Asymptotic Running Time of

Program Fragments

Tips for **code expert**

Outro



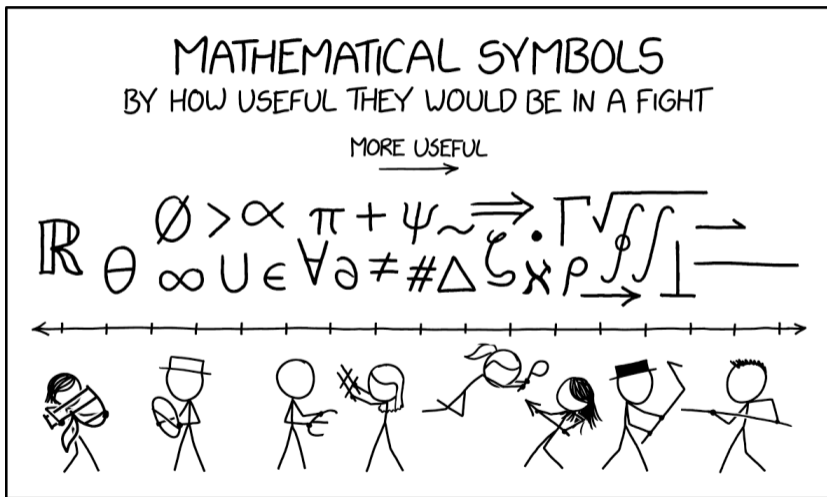
`n.ethz.ch/~agavranovic`

▶ Exercise Session Material

▶ Adel's Webpage

▶ Mail to Adel

Comic of the Week



1. Intro

Intro

Intro

- Who am I?

Intro

- Who am I?
- Who are you?

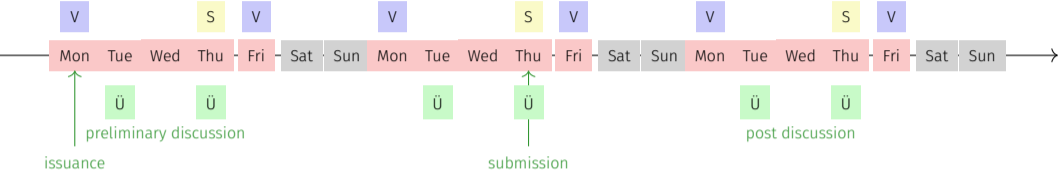
2. Learning Objectives

Learning Objectives

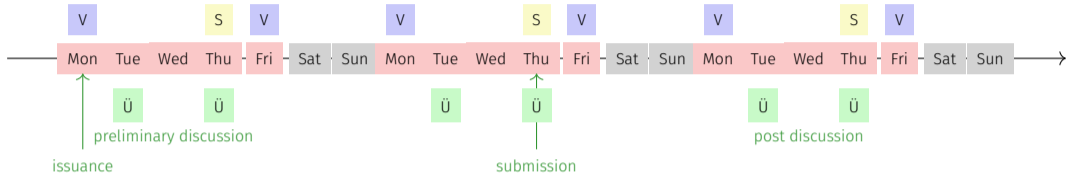
$O(n)$ $\mathcal{O}(n)$

- Get to know the weekly agenda for this course
- Understand differences between Problem, Algorithm, and Program
- Get to know big- \mathcal{O} notation (and its friends Ω and Θ)
- Learn some \LaTeX and Markdown

Exercises



Exercises



- Exercises available on Monday.
- Preliminary discussion in the following recitation session
- Solution of the exercise until the following Thursday.
- Discussion of the exercise in the next recitation session.
- Feedback roughly within 10 days after submission date.
- Study Center on Thursday.

4. Repetition Theory

Warm-up

- What is a problem?

Warm-up

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- What is an algorithm?

Warm-up

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- What is an algorithm?
 - well-defined computing procedure to compute output data from input data.

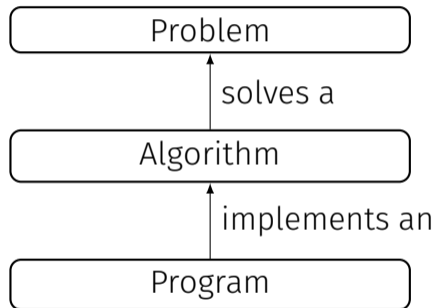
Warm-up

- What is a problem?
- What is an algorithm?
 - well-defined computing procedure to compute output data from input data.
- What is a program?

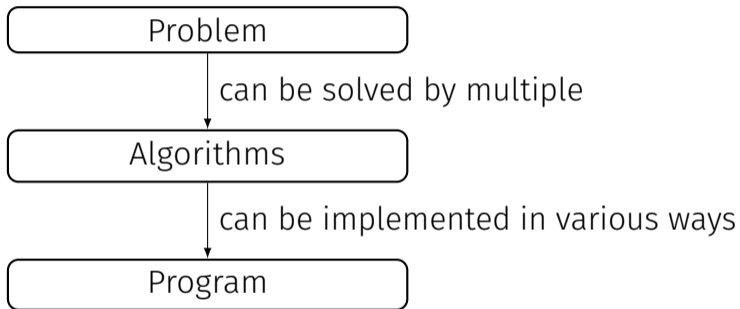
Warm-up

- What is a problem?
- What is an algorithm?
 - well-defined computing procedure to compute output data from input data.
- What is a program?
 - Concrete implementation of an algorithm

Problems, Algorithms and Programs



Warm-up



Efficiency

Program	Computing time	Measurable value on an actual machine.
Algorithm	<u>Cost</u>	Number of elementary operations
Problem	Complexity	Minimal (asymptotic) cost over all algorithms that solve the problem.

Efficiency

Program	Computing time	Measurable value on an actual machine.
Algorithm	Cost	Number of elementary operations
Problem	Complexity	Minimal (asymptotic) cost over all algorithms that solve the problem.

→ Estimation of cost or computing time depending on the input size, denoted by n .

Asymptotic behavior

- What are $\Omega(g(n))$, $\Theta(g(n))$, $\mathcal{O}(g(n))$?

Asymptotic behavior

- What are $\Omega(g(n))$, $\Theta(g(n))$, $\mathcal{O}(g(n))$?
- Sets of functions!

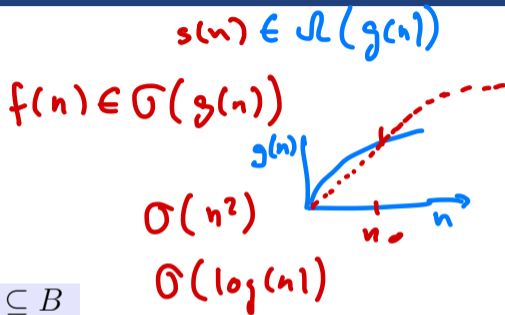
Asymptotic behavior

$n \rightarrow$

- What are $\Omega(g(n))$, $\Theta(g(n))$, $\mathcal{O}(g(n))$?
→ Sets of functions!

subset	$A \subseteq B$
proper subset	$A \subsetneq B$
intersection	$A \cap B$

σ



Asymptotic behavior

Given: function $f : \mathbb{N} \rightarrow \mathbb{R}$.

Definition:

$$\mathcal{O}(g) = \{f : \mathbb{N} \rightarrow \mathbb{R} \mid \exists c > 0, n_0 \in \mathbb{N} \mid \forall n \geq n_0 : 0 \leq f(n) \leq c \cdot g(n)\}$$

$$\Omega(g) = \{f : \mathbb{N} \rightarrow \mathbb{R} \mid \exists c > 0, n_0 \in \mathbb{N} \mid \forall n \geq n_0 : 0 \leq c \cdot g(n) \leq f(n)\}$$

$$\Theta(g) = \mathcal{O}(g) \cap \Omega(g)$$

Intuition:

$f \in \mathcal{O}(g)$: f grows asymptotically not faster than g . Algorithm with running time f is not worse than any other algorithm with g .

$f \in \Omega(g)$: f grows asymptotically not slower than g . Algorithm with running time f is worse than any other algorithm with g .

$f \in \Theta(g)$: f grows asymptotically as fast as g . Algorithm with running time f is as good as any other algorithm with g .

Used less often

Given: function $f : \mathbb{N} \rightarrow \mathbb{R}$.

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$$o(g) = \{f : \mathbb{N} \rightarrow \mathbb{R} \mid \forall c > 0 \exists n_0 \in \mathbb{N} \mid \forall n \geq n_0 : 0 \leq f(n) \leq c \cdot g(n)\}$$

$$\Omega(g) = \{f : \mathbb{N} \rightarrow \mathbb{R} \mid \exists c > 0, n_0 \in \mathbb{N} \mid \forall n \geq n_0 : 0 \leq c \cdot g(n) \leq f(n)\}$$

$$\omega(g) = \{f : \mathbb{N} \rightarrow \mathbb{R} \mid \forall c > 0 \exists n_0 \in \mathbb{N} \mid \forall n \geq n_0 : 0 \leq c \cdot g(n) \leq f(n)\}$$

$f \in o(g)$: f grows much slower than g

$f \in \omega(g)$: f grows much faster than g

Useful information for the exercise

Theorem 1

1. $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \Rightarrow f \in \mathcal{O}(g), \mathcal{O}(f) \subsetneq \mathcal{O}(g).$
2. $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = C > 0$ (C constant) $\Rightarrow f \in \Theta(g).$
3. $\frac{f(n)}{g(n)} \xrightarrow[n \rightarrow \infty]{} \infty \Rightarrow g \in \mathcal{O}(f), \mathcal{O}(g) \subsetneq \mathcal{O}(f).$

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Example 2

1. $\lim_{n \rightarrow \infty} \frac{n}{n^2} = 0 \Rightarrow n \in \mathcal{O}(n^2), \mathcal{O}(n) \subsetneq \mathcal{O}(n^2).$
2. $\lim_{n \rightarrow \infty} \frac{2n}{n} = 2 > 0 \Rightarrow 2n \in \Theta(n).$
3. $\frac{n^2}{n} \xrightarrow[n \rightarrow \infty]{} \infty \Rightarrow n \in \mathcal{O}(n^2), \mathcal{O}(n) \subsetneq \mathcal{O}(n^2).$

Property

$$f_1 \in \mathcal{O}(g), f_2 \in \mathcal{O}(g) \Rightarrow f_1 + f_2 \in \mathcal{O}(g)$$

4.1 Examples (Theory)

Examples

$$\mathcal{O}(g) = \{f : \mathbb{N} \rightarrow \mathbb{R} \mid \exists c > 0, \exists n_0 \in \mathbb{N} : \forall n \geq n_0 : 0 \leq f(n) \leq c \cdot g(n)\}$$

$f(n)$	$f \in \mathcal{O}(?)$	Example
$3n + 4$		
$2n$		
$n^2 + 100n$		
$n + \sqrt{n}$		

Examples

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$2n$	$\mathcal{O}(n)$	$c = 2, n_0 = 0$
$n^2 + 100n$	$\mathcal{O}(n^2)$	$c = 2, n_0 = 100$
$n + \sqrt{n}$	$\mathcal{O}(n)$	$c = 2, n_0 = 1$

Examples

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- $\mathcal{O}(n) \subseteq \mathcal{O}(n^2)$ is correct
- $\Theta(n) \subseteq \Theta(n^2)$ is wrong $n \notin \Omega(n^2) \supset \Theta(n^2)$

Quiz

$1 \in \mathcal{O}(15)$?

Quiz

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✓ better $1 \in \mathcal{O}(1)$

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$1 \in \mathcal{O}(15)$? ✓ better $1 \in \mathcal{O}(1)$

$2n + 1 \in \Theta(n)$?

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$n \in \Omega(\sqrt{n})$?

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A good strategy?

... Then I simply buy a new machine!

A good strategy?

... Then I simply buy a new machine! If today I can solve a problem of size n , then with a 10 or 100 times faster machine I can solve ...

Komplexität	(speed $\times 10$)	(speed $\times 100$)
$\log_2 n$		
n		
n^2		
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A good strategy?



... Then I simply buy a new machine! If today I can solve a problem of size n , then with a 10 or 100 times faster machine I can solve ... ¹

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n	$n \rightarrow 10 \cdot n$	$n \rightarrow 100 \cdot n$
n^2	$n \rightarrow 3.16 \cdot n$	$n \rightarrow 10 \cdot n$
2^n	$n \rightarrow n + 3.32$	$n \rightarrow n + 6.64$

¹To see this, you set $f(n') = c \cdot f(n)$ ($c = 10$ or $c = 100$) and solve for n'

4.2 Asymptotic Running Time of Program Fragments

Asymptotic Running Times with Θ

```
void run(int n){  
    for (int i = 1; i<n; ++i)  
        op();  
}
```

How often is `op()` called as a function of n ?

Asymptotic Running Times with Θ

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$$\sum_{i=1}^{n-1} 1 = n - 1 \in \Theta(n)$$

Asymptotic Running Times with Θ

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How often is `op()` called as a function of n ?

$$\sum_{i=1}^{n-1} \sum_{j=1}^{n-1} 1 = \sum_{i=1}^{n-1} (n-1) = (n-1) \cdot (n-1) \in \Theta(n^2)$$

Asymptotic Running Times with Θ

```
void run(int n){  
    for (int i = 1; i<n; ++i)  
        for (int j = i; j<n; ++j)  
            op();  
}
```

How often is `op()` called as a function of n ?

Asymptotic Running Times with Θ

```
void run(int n){  
    for (int i = 1; i<n; ++i)  
        for (int j = i; j<n; ++j)  
            op();  
}
```

How often is `op()` called as a function of n ?

$$\sum_{i=1}^{n-1} \sum_{j=i}^{n-1} 1 = \sum_{i=1}^{n-1} (n - i) = \sum_{i=1}^{n-1} i = \frac{n(n-1)}{2} \in \Theta(n^2)$$

Asymptotic Running Times with Θ

```
void run(int n){  
    for (int i = 1; i<n; ++i){  
        op();  
        for (int j = i; j<n; ++j)  
            op();  
    }  
}
```

How often is op() called?

Asymptotic Running Times with Θ

```
void run(int n){  
    for (int i = 1; i<n; ++i){  
        op();  
        for (int j = i; j<n; ++j)  
            op();  
    }  
}
```

How often is `op()` called?

$$\sum_{i=1}^{n-1} \left(1 + \sum_{j=i}^{n-1} 1 \right) = \sum_{i=1}^{n-1} (1 + (n - i)) = \underbrace{n - 1} + \underbrace{\frac{n(n - 1)}{2}} \in \Theta(n^2)$$

Asymptotic Running Times with Θ

```
void run(int n){  
    for (int i = 1; i<n; ++i){  
        op();  
        for (int j = 1; j<i*i; ++j)  
            op();  
    }  
}
```

How often is `op()` called?

Asymptotic Running Times with Θ

```
void run(int n){  
    for (int i = 1; i<n; ++i){  
        op();  
        for (int j = 1; j<i*i; ++j)  
            op();  
    }  
}
```

How often is op() called?

$$\sum_{i=1}^{n-1} \left(1 + \sum_{j=1}^{i^2-1} 1 \right) = \sum_{i=1}^{n-1} (1 + i^2 - 1) = \sum_{i=1}^{n-1} i^2 \in \Theta(n^3)$$

Asymptotic Running Times with Θ

```
void run(int n){  
    for(int i = 1; i <= n; ++i)  
        for(int j = 1; j*j <= n; ++j)  
            for(int k = n; k >= 2; --k)  
                op();  
}
```

Handwritten red annotations illustrating the complexity of the nested loops: a large bracket on the left is labeled n , a smaller bracket inside it is labeled \sqrt{n} , and a third bracket on the right is labeled n .

How often is `op()` called as a function of n ?

Asymptotic Running Times with Θ

```
void run(int n){  
    for(int i = 1; i <= n; ++i)  
        for(int j = 1; j*j <= n; ++j)  
            for(int k = n; k >= 2; --k)  
                op();  
}
```

$$\sqrt{n} \leftarrow j^2 = n$$

How often is `op()` called as a function of n ?

$$\sum_{i=1}^n \sum_{j=1}^{\lfloor \sqrt{n} \rfloor} n - 1 \in \Theta\left(\sum_{i=1}^n n^{3/2}\right) = \Theta(\sqrt{n^5})$$

Asymptotic Running Times with Θ

```
int f(int n){
```

```
  i=1;
```

```
  while (i <=  $\frac{n}{i}$ ){
```

```
    i = i*2;
```

```
    op();
```

```
  }
```

```
  return i;
```

```
}
```

$\log_2(n)$

$$\log_2(n^3) = 3 \log_2(n) \in \Theta(\log(n))$$

How often is op() called as a function of n ?

Asymptotic Running Times with Θ

```
int f(int n){
    i=1;
    while (i <= n*n*n){
        i = i*2;
        op();
    }
    return i;
}
```

How often is `op()` called as a function of n ?

$$|\{i \in \mathbb{N} : 2^i \leq n^3\}| \in \Theta(\log_2 n^3) = \Theta(\log n)$$

5. Appendix

Some formulas with derivation

Sums

$$\sum_{i=0}^n i = ?$$

Sums

$$\sum_{i=0}^n i = \frac{n \cdot (n + 1)}{2}$$

Sums

$$\sum_{i=0}^n i = \frac{n \cdot (n + 1)}{2}$$

Why?

Sums

$$\sum_{i=0}^n i = \frac{n \cdot (n + 1)}{2}$$

Why?

Intuition

$$1 + \dots + 100 = (1 + 100) + (2 + 99) + (3 + 98) + \dots + (50 + 51)$$

Sums

$$\sum_{i=0}^n i = \frac{n \cdot (n + 1)}{2}$$

Why?

Intuition

$$1 + \dots + 100 = (1 + 100) + (2 + 99) + (3 + 98) + \dots + (50 + 51)$$

More formally?

Sums

$$\sum_{i=0}^n (n - i) = ?$$

Sums

$$\sum_{i=0}^n (n - i) = \sum_{i=0}^n i$$

Sums

$$\sum_{i=0}^n (n - i) = \sum_{i=0}^n i$$

$$\begin{aligned} \Rightarrow 2 \cdot \sum_{i=0}^n i &= \sum_{i=0}^n i + \sum_{i=0}^n (n - i) \\ &= \sum_{i=0}^n (i + (n - i)) = \sum_{i=0}^n n = (n + 1) \cdot n \end{aligned}$$

Sums

$$\sum_{i=0}^n (n - i) = \sum_{i=0}^n i$$

$$\begin{aligned} \Rightarrow 2 \cdot \sum_{i=0}^n i &= \sum_{i=0}^n i + \sum_{i=0}^n (n - i) \\ &= \sum_{i=0}^n (i + (n - i)) = \sum_{i=0}^n n = (n + 1) \cdot n \end{aligned}$$

Sums

$$\sum_{i=0}^n i^2 = ?$$

Sums

$$\sum_{i=0}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

Sums

$$\sum_{i=0}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

This you do not need to know by heart. But you should ~~know~~
~~that it is a polynomial of third degree.~~

show it on your CS.

Sums

How do you derive something like this?

Sums

How do you derive something like this? Interesting Trick: On the one hand

$$\sum_{i=0}^n i^3 - \sum_{i=1}^n (i-1)^3 = \sum_{i=0}^n i^3 - \sum_{i=0}^{n-1} i^3 = n^3,$$

Sums

How do you derive something like this? Interesting Trick: On the one hand

$$\sum_{i=0}^n i^3 - \sum_{i=1}^n (i-1)^3 = \sum_{i=0}^n i^3 - \sum_{i=0}^{n-1} i^3 = n^3,$$

on the other hand

$$\begin{aligned} \sum_{i=0}^n i^3 - \sum_{i=1}^n (i-1)^3 &= \sum_{i=1}^n i^3 - \sum_{i=1}^n (i-1)^3 \\ &= \sum_{i=1}^n i^3 - (i-1)^3 = \sum_{i=1}^n 3 \cdot i^2 - 3 \cdot i + 1 \end{aligned}$$

Exponents and Logarithms

$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = ?$$

$$\frac{a^x}{a^y} = ?$$

$$a^{x \cdot y} = ?$$

$$\log_b x = ?$$

$$\log_a (x \cdot y) = ?$$

$$\log_a \frac{x}{y} = ?$$

$$\log_a x^y = ?$$

$$\log_a n! = ?$$

$$a^{\log_b x} = ?$$

Exponents and Logarithms

$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = a^{x+y}$$

$$\frac{a^x}{a^y} = ?$$

$$a^{x \cdot y} = ?$$

$$\log_b x = ?$$

$$\log_a(x \cdot y) = ?$$

$$\log_a \frac{x}{y} = ?$$

$$\log_a x^y = ?$$

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$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = a^{x+y}$$

$$\frac{a^x}{a^y} = a^{x-y}$$

$$a^{x \cdot y} = ?$$

$$\log_b x = ?$$

$$\log_a(x \cdot y) = ?$$

$$\log_a \frac{x}{y} = ?$$

$$\log_a x^y = ?$$

$$\log_a n! = ?$$

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$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

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$$\log_b x = ?$$

$$\log_a (x \cdot y) = \log_a x + \log_a y$$

$$\log_a \frac{x}{y} = ?$$

$$\log_a x^y = ?$$

$$\log_a n! = ?$$

$$a^{\log_b x} = ?$$

Exponents and Logarithms

$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = a^{x+y}$$

$$\frac{a^x}{a^y} = a^{x-y}$$

$$a^{x \cdot y} = (a^x)^y$$

$$\log_b x = ?$$

$$\log_a(x \cdot y) = \log_a x + \log_a y$$

$$\log_a \frac{x}{y} = \log_a x - \log_a y$$

$$\log_a x^y = ?$$

$$\log_a n! = ?$$

$$a^{\log_b x} = ?$$

Exponents and Logarithms

$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = a^{x+y}$$

$$\frac{a^x}{a^y} = a^{x-y}$$

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$$\log_a (x \cdot y) = \log_a x + \log_a y$$

$$\log_a \frac{x}{y} = \log_a x - \log_a y$$

$$\log_a x^y = y \log_a x$$

$$\log_a n! = ?$$

$$a^{\log_b x} = ?$$

Exponents and Logarithms

$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = a^{x+y}$$

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$$\log_a(x \cdot y) = \log_a x + \log_a y$$

$$\log_a \frac{x}{y} = \log_a x - \log_a y$$

$$\log_a x^y = y \log_a x$$

$$\log_a n! = \sum_{i=1}^n \log i$$

$$\log_b x = ?$$

$$a^{\log_b x} = ?$$

Exponents and Logarithms

$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = a^{x+y}$$

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$$\log_b x = \log_b a \cdot \log_a x$$

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$$\log_a \frac{x}{y} = \log_a x - \log_a y$$

$$\log_a x^y = y \log_a x$$

$$\log_a n! = \sum_{i=1}^n \log i$$

$$a^{\log_b x} = ?$$

Exponents and Logarithms

$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = a^{x+y}$$

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$$a^{\log_b x} = x^{\log_b a}$$

Exponents and Logarithms

$$\log_a y = x \Leftrightarrow a^x = y \quad (a > 0, y > 0)$$

$$a^x \cdot a^y = a^{x+y}$$

$$\frac{a^x}{a^y} = a^{x-y}$$

$$a^{x \cdot y} = (a^x)^y$$

$$\log_a(x \cdot y) = \log_a x + \log_a y$$

$$\log_a \frac{x}{y} = \log_a x - \log_a y$$

$$\log_a x^y = y \log_a x$$

$$\log_a n! = \sum_{i=1}^n \log i$$

$$a^{\log_b x} = x^{\log_b a}$$

$$\log_b x = \log_b a \cdot \log_a x$$

To see the last line, replace $x \rightarrow a^{\log_a x}$

Comparisons

$$\frac{n^2}{2^n} \xrightarrow{n \rightarrow \infty} ?$$

Comparisons

$$\frac{n^2}{2^n} \xrightarrow{n \rightarrow \infty} 0$$

Comparisons

$$\frac{n^{10000}}{2^n} \xrightarrow{n \rightarrow \infty} ?$$

Comparisons

$$\frac{n^{10000}}{2^n} \xrightarrow{n \rightarrow \infty} 0$$

Comparisons

$$d > 1, c > 0$$

$$\frac{n^c}{d^n} \xrightarrow{n \rightarrow \infty} ?$$

Comparisons

$$d > 1, c > 0$$

$$\frac{n^c}{d^n} \xrightarrow{n \rightarrow \infty} 0$$

Comparisons

$$d > 1, c > 0$$

$$\frac{n^c}{d^n} \xrightarrow{n \rightarrow \infty} 0$$

because

$$\frac{n^c}{d^n} = \frac{2^{\log_2 n^c}}{2^{\log_2 d^n}} = 2^{c \cdot \log_2 n - n \log_2 d}$$

Comparisons

$$\frac{n}{\log n} \xrightarrow{n \rightarrow \infty} ?$$

Comparisons

$$\frac{n}{\log n} \xrightarrow{n \rightarrow \infty} \infty$$

Comparisons

$$\frac{n \log n}{\sqrt{n}} \xrightarrow{n \rightarrow \infty} ?$$

Comparisons

$$\frac{n \log n}{\sqrt{n}} \xrightarrow{n \rightarrow \infty} \infty$$

Comparisons

$$\frac{\log_2 n^2}{\sqrt{n}} \xrightarrow{n \rightarrow \infty} ?$$

Handwritten in red: = $2 \log_2(n)$

Comparisons

$$\frac{\log_2 n^2}{\sqrt{n}} \xrightarrow{n \rightarrow \infty} 0$$

Comparisons

$$\frac{\log_2 n^2}{\sqrt{n}} \xrightarrow{n \rightarrow \infty} 0$$

$$\log_2 n^2 = 2 \log_2 n$$

$$\sqrt{n} = n^{1/2} = 2^{\log_2 n^{1/2}} = (\sqrt{2})^{\log_2 n}$$

$$\frac{\log n^2}{\sqrt{n}} = 2 \frac{\log_2 n}{(\sqrt{2})^{\log_2 n}}$$

which behaves because of $\log_2 n \rightarrow \infty$ for $n \rightarrow \infty$ like $2 \frac{n}{(\sqrt{2})^n}$

6. Tips for **code** expert

Tips for **code expert** Exercise 1

All Text Tasks

Tips for **code expert** Exercise 1

All Text Tasks

- Please learn a little \LaTeX and Markdown. It will make your (and my) life a lot easier

Tips for **code expert** Exercise 1

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- Please learn a little \LaTeX and Markdown. It will make your (and my) life a lot easier
- Useful Links and some tools I use

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- [▶ Mathpix Snipping Tool \(paid\)](#)

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Task "Some Proofs"

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Task "Some Proofs"

- No need for a rigorous proof (this is not Disk Math)

Tips for **code expert** Exercise 1

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Task "Some Proofs"

- No need for a rigorous proof (this is not Disk Math)
- It pays off to revisit some of the log-properties that we've covered today

Tips for **code expert** Exercise 2

Task "Prefix Sum in 2D"

Tips for **code expert** Exercise 2

Task "Prefix Sum in 2D"

- Study the Prefix Sum in 1D well and go from there
- Make sketches!

Tips for **code expert** Exercise 2

Task "Prefix Sum in 2D"

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Task "Sliding Window"

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Task "Proofs by Induction"

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Task "Sliding Window"

- Sketches!

Task "Proofs by Induction"

- The binomial formula will be useful for the second one

Tips for **code expert** Exercise 2

Task "Prefix Sum in 2D"

- Study the Prefix Sum in 1D well and go from there
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Task "Sliding Window"

- Sketches!

Task "Proofs by Induction"

- The binomial formula will be useful for the second one

Task "Karatsuba Ofman"

Tips for **code expert** Exercise 2

Task "Prefix Sum in 2D"

- Study the Prefix Sum in 1D well and go from there
- Make sketches!

Task "Sliding Window"

- Sketches!

Task "Proofs by Induction"

- The binomial formula will be useful for the second one

Task "Karatsuba Ofman"

- Just translate the math into code

7. Outro

General Questions?

See you next time

Have a nice week!