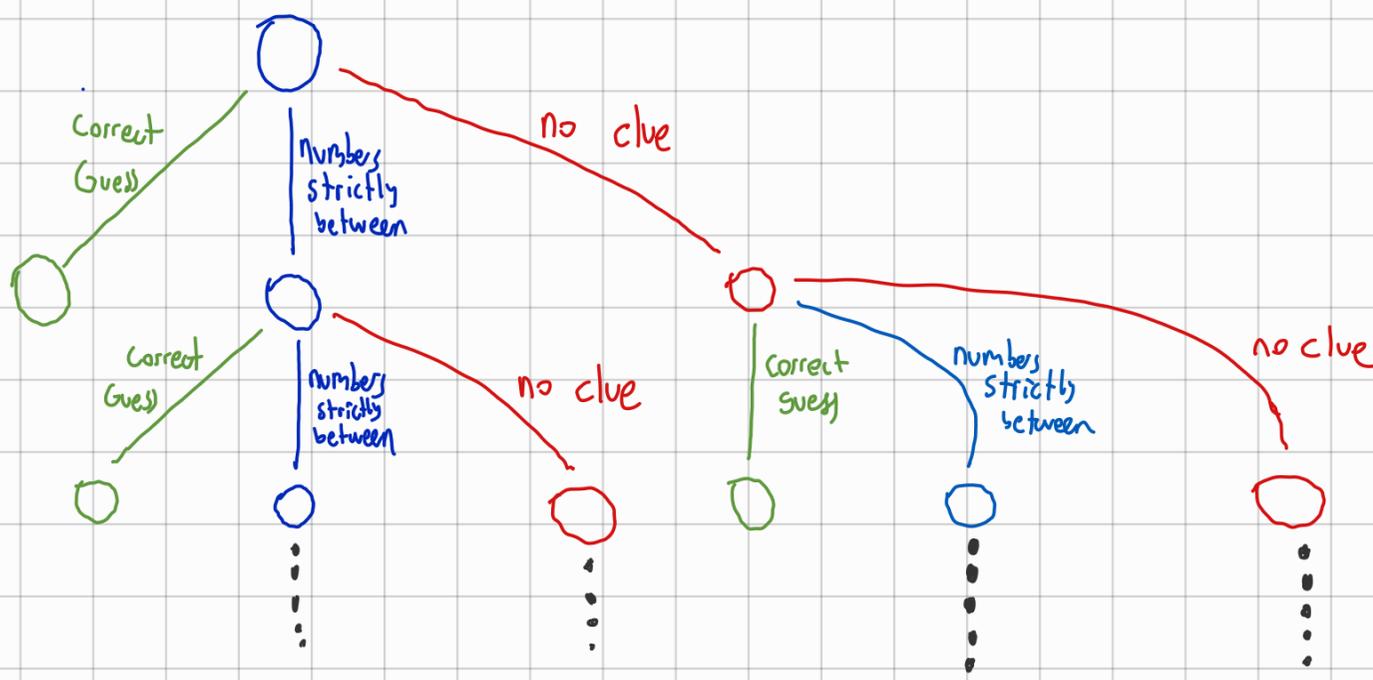


# EXERCISE 5.2

After each guess :

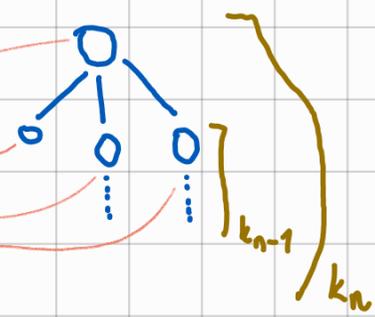
- i) Bob guesses correctly
- ii) Alice tells numbers strictly between
- iii) Alice doesn't give any clue



$k_n$  = maximum number of vertices in a tree of depth  $n \in \mathbb{N}_0$ .

$k_0 = 1$      $k_1 = 4$      $k_2 = 10$

$k_0 = 1, \quad \forall n \geq 1 \quad k_n = 2k_{n-1} + 1 + 1$



Find a non-recursive formula for  $k_n$ .  
 First guess it as  $3 \cdot 2^n - 2$ , then prove by induction.

Total numbers of miss Alice can choose :  $\lfloor 200 \rfloor - \frac{200 \cdot 199}{2} = 19900$

The tree representing Bob's strategy must have at least 19900 leaves.

Note that the depth of the tree is the number of guesses Bob needs to make in the worst case.

$$k_{12} = 12286$$

Bob cannot certainly win in at most 12 attempts.

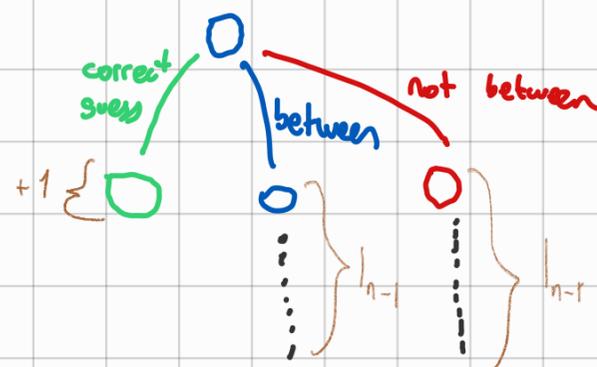
We need 19900 leaves but the decision tree with depth 12 has  $< 12286$  leaves.

b) If the decision tree with depth 14 has at least 19900 leaves, then Bob can certainly win in at most 14 attempts. Let's determine the number of leaves.

$f_n$ : maximum number of leaves in a tree of depth  $n$ .

$$f_1 = 1, \forall n > 1 \quad f_n = 2f_{n-1} + 1$$

Guess  $f_n = 2^n - 1$ , then prove by induction.



So, the tree with depth 14 has  $2^{14} - 1 = 16383$  leaves. Hence Bob cannot certainly win in at most 14 attempts.

## EXERCISE 5.3

a) After each iteration of the while loop, the depth of  $N$  gets decreased by 1. Since  $H$  is a complete binary tree, it has at most  $O(\log n)$  depth. Hence  $N$  can be pushed  $O(\log n)$ -times up.

b)

i) Algorithm terminated because  $N_{stop}$  was the root node. Heap condition satisfied because  $N_{stop}$  is the root node.

ii) Algorithm terminated because  $key(P) \geq key(N_{stop})$ . So  $N_{stop}$  satisfies the heap condition. Besides, all nodes with depth less than or equal to depth of  $N_{stop}$  have their keys unchanged. So these keys still satisfy heap condition.

c) Induction on  $t$ , number of iterations of while loop. Base case is  $t=0$ .

## EXERCISE 5.4

a) Linked List. Push & Pop requires  $O(1)$  time.

b) Doubly Linked List. Enqueue & Dequeue requires  $O(1)$  time.

