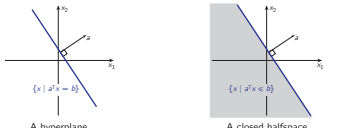
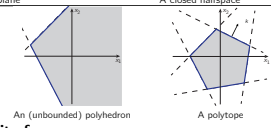
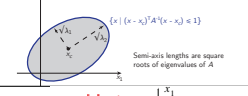
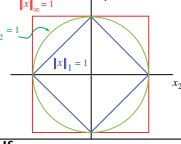


MPC Summary	
Jorit Geurts - jgeurts@ethz.ch Version: 17. August 2023	
1	Systems Theory
1.1	System Dynamics
1.1.1	Continuous Time
Nonlinear Time-Invariant Continuous Time State Space $\dot{x} = g(x, u) \quad x \in \mathbb{R}^n, u \in \mathbb{R}^m \quad g: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}^n$ $y = h(x, u) \quad y \in \mathbb{R}^p \quad h: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}^p$	
LTI Continuous Time State Space Linearization using Taylor Expansion around operating point: $f(x) \approx f(\bar{x}) + \left. \frac{\partial f}{\partial x} \right _{\bar{x}} (x - \bar{x})$	
Resulting system: $\left. \begin{aligned} \dot{x} &= \left. \frac{\partial g}{\partial x} \right _{x_s, u_s} \delta x + \left. \frac{\partial g}{\partial u} \right _{x_s, u_s} \delta u \\ y &= \left. \frac{\partial h}{\partial x} \right _{x_s, u_s} \delta x + \left. \frac{\partial h}{\partial u} \right _{x_s, u_s} \delta u \end{aligned} \right\} \begin{aligned} \dot{x} &= A^c x + B^c u \\ y &= C x + D u \end{aligned}$ $C \in \mathbb{R}^{p \times n} \quad D \in \mathbb{R}^{p \times m}$	
Solution: $e^{A^c t} = \sum_{n=0}^{\infty} \frac{(A^c t)^n}{n!}$ $x(t) = e^{A^c(t-t_0)} x_0 + \int_{t_0}^t e^{A^c(t-\tau)} B^c u(\tau) d\tau$	
1.1.2	Discrete Time
Euler Discretization (T_s = sampling time) (stability not guaranteed) $\dot{x}^c \approx \frac{x^c(t+T_s) - x^c(t)}{T_s}, x(k) := x^c(t_0 + kT_s), u(k) := u^c(t_0 + kT_s)$	
Nonlinear System: $x(k+1) = x(k) + T_s(g^c(x(k), u(k))) = g(x(k), u(k))$ $y(k) = h^c(x(k), u(k)) = h(x(k), u(k))$	
Linear System: $x(k+1) = A^d x(k) + B^d u(k), \quad A^d = \mathbb{I} + T_s A^c, \quad B^d = T_s B^c$ $y(k) = C^d x(k) + D^d u(k), \quad C^d = C^c, \quad D^d = D^c$	
Exact Discretization (only for linear systems), (stability guaranteed) Exact solution (u assumed constant over T_s): $x(t_{k+1}) = \underbrace{e^{A^c T_s}}_{=A} x(t_k) + \underbrace{\int_0^{T_s} e^{A^c(T_s-\tau)} B^c d\tau}_{=B(A-A^c)^{-1}(A-\mathbb{I})B^c} u(t_k)$	
We see the solution over k is then given by: $x(k+N) = A^N x(k) + \sum_{i=0}^{N-1} A^i B u(k+N-1-i)$	
1.2	Linear System Analysis
DT Stability (Lypapunov indirect method) $x(k+1) = Ax(k)$ stable iff $ \lambda_j < 1, \forall j \rightarrow$ NL system stable if $ \lambda_i = 1$ NL system no info, if $ \lambda_i > 1$ NL system unstable	
LTI DT Controllability can reach x^* from $x(0)$ in n steps $C = [B \quad \dots \quad A^{n-1}B] \Rightarrow \text{rank}(C) \stackrel{!}{=} n$	
DT Observability uniquely distinguish IC from output $\mathcal{O} = [C^T \quad \dots \quad (CA^{n-1})^T]^T \Rightarrow \text{rank}(\mathcal{O}) \stackrel{!}{=} n$	
Stabilizability iff all uncontrollable modes stable $\Lambda_A^+ = \{\lambda 1 \leq \lambda \}$ if $\text{rank}(\{\lambda_j \mathbb{I} - A \mid B\}) = n \quad \forall \lambda_j \in \Lambda_A^+ \Rightarrow (A, B)$ stabilizable	
Detectability iff all unobservable modes stable $\Lambda_A^+ = \{\lambda 1 \leq \lambda \}$ if $\text{rank}([A^T - \lambda_j \mathbb{I} \mid C^T]) = n \quad \forall \lambda_j \in \Lambda_A^+ \Rightarrow (A, C)$ detect.	
1.3	Nonlinear System Analysis
Lypapunov Stability (w.r.t eq. point \bar{x} of a system) Lypapunov Stable if for every $\epsilon > 0$ exists $\delta(\epsilon)$ s.t. $\ x(0) - \bar{x}\ < \delta(\epsilon) \Rightarrow \ x(k) - \bar{x}\ < \epsilon$ Globally Asympt. Stable if Lyp. stable & Attractive $\lim_{k \rightarrow \infty} \ x(k) - \bar{x}\ = 0 \quad \forall x(0) \in \Omega$	
Lypapunov Function Consider eq point $\bar{x} \in \mathbb{R}^n$. $V: \mathbb{R}^n \rightarrow \mathbb{R}$, continuous at origin, finite $\forall x$, (1) $\ x\ \rightarrow 0 \Rightarrow V(x) \rightarrow \infty$ (2) $V(0) = 0, \quad V(x) > 0 \quad \forall x \in \mathbb{R}^n \setminus \{0\}$ (3) $V(g(x)) - V(x) < -\alpha(x) \quad \forall x \in \mathbb{R}^n$ where $\alpha: \mathbb{R}^n \rightarrow \mathbb{R}$ continuous pos. def.	
Lypapunov Stability If sys admits a $V(x) \Rightarrow x = 0$ is Globally Asympt. Stable Caution if α pos. semidef $\Rightarrow x = 0$ is Globally Lypapunov Stable Asympt. Stable in pos invar set $\Omega \subset \mathbb{R}^n$ if Lyp. stable and attactive $\lim_{k \rightarrow \infty} \ x(k) - \bar{x}\ = 0 \quad \forall x(0) \in \Omega$	
Globally Asympt. Stable if asympt. stable & $\Omega = \mathbb{R}^n$	

2	Linear Quadratic Optimal Control
2.1	Linear Quadratic Optimal Control
Problem Definition $J(x(0), U) := x_N^T P x_N + \sum_{i=0}^{N-1} (x_i^T Q x_i + u_i^T R u_i)$ subj. to $x_{i+1} = A x_i + B u_i, \quad x_0 = x(0)$ • N : horizon length • $Q \succeq 0, \quad Q = Q^T$ • $x(0)$: current state • $P \succeq 0, \quad P = P^T$ • $R \succ 0, \quad R = R^T$ • x_i, u_i : opt. variable	
2.1.1	Batch Approach
Idea explicitly represent $x_i \in \mathbb{R}^n$ through x_0 & $u_i \in \mathbb{R}^m$ $\begin{bmatrix} x_0 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} \mathbb{I} \\ A \\ \vdots \\ A^N \end{bmatrix} x(0) + \begin{bmatrix} 0 & \dots & 0 \\ B & \dots & 0 \\ \vdots & \ddots & \vdots \\ A^{N-1}B & \dots & B \end{bmatrix} \begin{bmatrix} u_0 \\ \vdots \\ u_{N-1} \end{bmatrix}$	
Equivalent to $S^x \in \mathbb{R}^{(N+1)n \times n}, \quad S^u \in \mathbb{R}^{(N+1)n \times Nm}$ $X = S^x x(0) + S^u U \rightarrow J(x(0), U) = X^T \bar{Q} X + U^T \bar{R} U$ Cost: $\bar{Q} := \text{blockdiag}(Q, \dots, Q, P)$ & $\bar{R} := \text{blockdiag}(R, \dots, R)$ Solve by setting gradient to zero: $2HU^* + F^T x(0) = 0$ Optimal Input: $H = (S^u)^T \bar{Q} S^u + \bar{R}, F = (S^x)^T \bar{Q} S^u$	
$U^*(x(0)) = -((S^u)^T \bar{Q} S^u + \bar{R})^{-1} (S^u)^T \bar{Q} S^x x(0)$	
Optimal Cost $J^* = x(0)^T [S^x \bar{Q} S^x - S^x \bar{Q} S^u (S_u^T \bar{Q} S_u + \bar{R})^{-1} S_u^T \bar{Q} S^x x(0)]$	
2.1.2	Recursive Approach
Idea: Recursively compute optimal input u_j^* and optimal cost J_j^* $J_j^*(x(j)) := \min_{U_j \rightarrow N} x_N^T P x_N + \sum_{i=j}^{N-1} (x_i^T Q x_i + u_i^T R u_i)$ $P = P_N, F \leftarrow f(P),$ Control input, $P \leftarrow f(F),$ Cost calculation, repeat Optimal Control Policy $u_i^* = -(B^T P_{i+1} B + R)^{-1} B^T P_{i+1} A \cdot x(i) := F_i x_i$	
Optimal Cost-To-Go $J_i^*(x_i) = x_i^T P_i x_i$ RDE – Riccati Difference Equation ($P_N = P$) $P_i = A^T P_{i+1} A + Q - A^T P_{i+1} B (B^T P_{i+1} B + R)^{-1} B^T P_{i+1} A$ Numerically Safer Alternative $P_i = Q + F_i^T R F_i + (A + B F_i)^T P (A + B F_i)$	
2.1.3	Comparison - Batch vs. Recursive
• Batch – sequence of numeric values U^* • Recursive – feedback policies u_i^* • Control actions identical if perfect model • Disturbances – Recursive more robust to disturbances • Computational efficiency – Recursive more efficient for large N – Matrix inversion in Batch approach expensive • Constraints – Neither works with constraints on x_i or u_i • Batch Approach easier to adapt when constraints are present constrained minimization (solving for J_{i+1} with constraints) hard	
2.2	Receding Horizon Control
Compute optimal control policy for N steps, apply only first step, then re-compute $U^* := \text{argmin} x_N^T P x_N + \sum_{i=0}^{N-1} x_i^T Q x_i + u_i^T R u_i$ subj. to $x_{i+1} = A x_i + B u_i \Rightarrow U^*$	
• Extract first input in sequence: $U^* = \{u_0^*, \dots, u_{N-1}^*\} \Rightarrow u_0^*$ • Introduce feedback to sys: $x(k+1) = Ax(k) + Bu(k) \Rightarrow x$ Why Reoptimize Provides robustness to noise / modeling errors, Sol'n at k subopt. (finite horizon) \rightsquigarrow reopt. potentially better performance	
2.3	Infinite Horizon Control LQR
Solve LQOC for $N \rightarrow \infty$ $J_{\infty}(x(0)) = \min_u \sum_{i=0}^{\infty} x_i^T Q x_i + u_i^T R u_i$ subj. to $x_{i+1} = A x_i + B u_i, \quad x_0 = x(0)$	
As with recursive approach it must hold: $u^*(k) = -(B^T P_{\infty} B + R)^{-1} B^T P_{\infty} A \cdot x(k) := F_{\infty} x(k)$	
with infinite cost to go: $J_{\infty}(x(k)) = (x(k))^T P_{\infty} x(k)$ Algebraic Riccati Equation (ARE) to find P_{∞} : $P_{\infty} = A^T P_{\infty} A + Q - A^T P_{\infty} B (B^T P_{\infty} B + R)^{-1} B^T P_{\infty} A$ LQR Lypapunov Function If (A, B) stabilizable, $(Q^{1/2}, A)$ detectable $\rightsquigarrow J^*(x) = x^T P_{\infty} x$ is Lyp. func. for system $x^+ = (A + B F_{\infty}) x$	
Choice of P in Finite Horizon Control • Can choose to match ∞ -Horizon sol'n \rightsquigarrow Make $P \approx J_{N \rightarrow \infty}$ with ARE • Can Choose P assuming no control action after end of horizon This P determined from solving Lyp eqn $A^T P A + Q = P$ Only makes sense if system asympt. stable • Assume we want state and input both to be 0 at end of horizon \rightsquigarrow no P but extra constraint $x_{i+N} = 0$	

3	Convex Sets
3.1	Problem Formulation
$\begin{aligned} \min_{x \in \text{dom}(f)} f(x) \quad & \text{subj. to } g_i(x) \leq 0 \quad i = 1, \dots, m \\ & h_i(x) = 0 \quad i = 1, \dots, p \end{aligned}$	
• $\mathcal{X} := \{x \in \text{dom}(f) \mid g_i \leq 0, h_i = 0\}$ feasible set • g_i : ineq constraints, h_i : eq constraints Feasibility Point x satisfies $g_i \leq 0, h_i = 0$ & eq constraints Optimal Value lowest cost $p^* = f(x^*) = \min_{x \in \mathcal{X}} f(x)$ Strictly Feasible Point x satisfies $g_i < 0$ Optimizer smallest $p^*, x \in \mathcal{X}$: $\text{argmin}_{x \in \mathcal{X}} f(x) := \{x \in \mathcal{X} \mid f(x) = p^*\}$ Caution NOT always unique Active Constraints: when ineq const. are eq \rightsquigarrow "active" Locally Optimal: $y \in \mathcal{X}, \ y - x\ \leq R \Rightarrow f(y) \geq f(x)$ Unbounded Below $p^* = -\infty$, Unconstrained $\mathcal{X} = \mathbb{R}^n$ Redundant Constraints do not change feasible set Globally Optimal: $y \in \mathcal{X} \Rightarrow f(y) \geq f(x)$ Infeasible $p^* = \infty \Leftrightarrow \mathcal{X} = \{\}$	
3.2	Convex Sets
Definition Set \mathcal{X} is convex iff for any pair of points x and y in \mathcal{X} : $\lambda x + (1 - \lambda)y \in \mathcal{X} \quad \forall \lambda \in [0, 1], \quad \forall x, y \in \mathcal{X}$ Interpretation: All lines starting in \mathcal{X} stay within \mathcal{X} Convex Combination: $x = \theta_1 x_1 + \theta_2 x_2 + \dots + \theta_k x_k, \quad \text{with } \sum_i \theta_i = 1, \theta_i \geq 0$	
Hyperplane $\{x \in \mathbb{R}^n \mid a^T x = b\}$ Halfspace $\{x \in \mathbb{R}^n \mid a^T x \leq b\}$ open: <, closed: \leq	
	
Polyhedron $P := \{x \mid a_i^T x \leq b_i, i = \dots\}$ $:= \{x \mid Ax \leq b\}$	
	
Polytope: bounded Polyhedron Intersection of Polytopes in inequality form: $\{x \mid Ax \leq b\} \cap \{x \mid Cx \leq d\} = \{x \mid \begin{bmatrix} A \\ C \end{bmatrix} x \leq \begin{bmatrix} b \\ d \end{bmatrix}\}$	
Ellipsoid $\{x \mid (x - x_c)^T A^{-1} (x - x_c) \leq 1\}$ x_c : center of ellipsoid 	
Norm Ball $\{x \mid \ x - x_c\ \leq r\}$	
• $p = 2$ Euclidean Norm $\ x\ _2 = \sqrt{\sum_i x_i^2}$ • $p = 1$ Sum of Absolute $\ x\ _1 = \sum_i x_i $ • $p = \infty$ Largest Absolute 	
Intersection \cap of two convex sets is convex itself Union \cup of two convex sets is NOT convex in general	
3.3	Convex Functions
Definition A function $f: \mathcal{D} \rightarrow \mathbb{R}$ is convex if and only if its domain $\mathcal{D} = \text{dom}(f)$ is a convex set and $f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y), \quad \forall x, y \in \mathcal{D}, \lambda \in [0, 1]$ $f: \mathcal{D} \rightarrow \mathbb{R}$ is strictly convex if the inequality is strict. f is concave if $-f$ is convex.	
First order condition A differentiable function $f: \mathcal{D} \rightarrow \mathbb{R}$ with a convex domain \mathcal{D} is convex if and only if $f(y) \geq f(x) + \nabla f(x)^T (y - x), \quad \forall x, y \in \mathcal{D}$	
Gradient is given by: $\nabla f(x) = \left[\frac{\partial f(x)}{\partial x_1}, \dots, \frac{\partial f(x)}{\partial x_n} \right]^T$	
Second order condition A twice-differentiable function $f: \mathcal{D} \rightarrow \mathbb{R}$ with a convex domain \mathcal{D} is convex if and only if $\nabla^2 f(x) \succeq 0, \quad \forall x \in \mathcal{D}, \quad \nabla^2 f(x)_{ij} = \frac{\partial^2 f(x)}{\partial x_i \partial x_j}$	
Strictly convex if $\nabla^2 f(x) \succ 0$. Examples Convex Affine $ax + b$ for any $a, b \in \mathbb{R}$ Exp. e^{ax} for any $A \in \mathbb{R}$ Powers $x^\alpha, x \in \mathbb{R}_{++}, \alpha \geq 1, \alpha \leq 0$ Vector norms on \mathbb{R}^n : $\ x\ _p = (\sum_{i=1}^n x_i ^p)^{1/p}, p \geq 1$	
Convexity Preserving Operations - Nonnegative weighted sum: $f(x) = \sum_{i=1}^n \theta_i f_i(x), \theta_i \geq 0$ - Composition with an affine mapping: $f(x) = g(Ax + b)$ - Pointwise maximum: $f(x) = \max\{f_1(x), \dots, f_m(x)\}$ - Partial minimization: $f(x, y) = \min_z g(x, y, z)$	
Concave Affine $ax + b$ for any $a, b \in \mathbb{R}$ Powers $x^\alpha, x \in \mathbb{R}_{++},$ for $0 \leq \alpha \leq 1$ Log $\log x$ on domain \mathbb{R}_{++} Entropy $-x \log x$ on domain \mathbb{R}_{++}	

3.3.1	Level and sublevel sets						
Levelset Definition The level set L_α of a function $f: \mathcal{D} \rightarrow \mathbb{R}$ is the set of points in the domain \mathcal{D} which $f(x) = \alpha$. $L_\alpha = \{x \in \mathcal{D} \mid f(x) = \alpha\}$ For $f: \mathbb{R}^2 \rightarrow \mathbb{R}$, these are contour lines of constant height.							
Sublevelset Definition The sublevel set C_α of a function $f: \mathcal{D} \rightarrow \mathbb{R}$ is the set of points in the domain \mathcal{D} which $f(x) \leq \alpha$. $C_\alpha = \{x \in \mathcal{D} \mid f(x) \leq \alpha\}$ If f is convex, then C_α is convex for all α .							
3.4	Convex Optimization Problem						
A convex optimization problem in standard form: $\min_{x \in \text{dom}(f)} f(x), \quad \text{subj. to } g_i(x) \leq 0 \quad i = 1, \dots, m$ $h_i(x) = a_i^T x = b_i \quad i = 1, \dots, p$							
f, g_i are convex, h_i are affine. Affine constraints are typically written in matrix form as $Ax = b$. Important Property: Feasible set of a convex optimization problem is convex. Local and Global Optimality: For a convex optimization problem, any local optimal solution is also a global optimal solution.							
Equivalent Optimization Problems Two problems are called equivalent if the solution from one can be inferred easily from the solution of the other. Example: $\min_x f(A_0 x + b)$ subj. to $g_i(A_i x + b_i) \leq 0, i = 1, \dots, m$ is equal to $\min_{x, y} f(y_0)$ subj. to $g_i(y_i) \leq 0, \quad A_i x + b_i = y_i, \quad i = 0, \dots, m$							
3.4.1	Linear Program						
<table border="0"> <tr> <th>Problem</th><th>Solutions</th></tr> <tr> <td>$\min_{x \in \mathbb{R}^n} c^T x$</td><td>Case 1: LP unbounded: $p^* = -\infty$</td></tr> <tr> <td></td><td>Case 2: Bounded and unique</td></tr> </table>		Problem	Solutions	$\min_{x \in \mathbb{R}^n} c^T x$	Case 1: LP unbounded: $p^* = -\infty$		Case 2: Bounded and unique
Problem	Solutions						
$\min_{x \in \mathbb{R}^n} c^T x$	Case 1: LP unbounded: $p^* = -\infty$						
	Case 2: Bounded and unique						
subj. to $Gx \leq h, \quad Ax = b$ Case 3: LP bounded but not unique							
3.4.2	Quadratic Program						
Problem \rightarrow solution is unique $\min_{x \in \mathbb{R}^n} \frac{1}{2} x^T H x + q^T x + r$ subj. to $Gx \leq h, \quad Ax = b$							
r not needed (does not change optimal x), Convex if $H \succ 0$ Case 1: optimizer lies strictly inside the feasible polyhedron Case 2: optimizer lies on the boundary of the feasible polyhedron							
3.5	Optimality Conditions						
3.5.1	Lagrang Dual Problem						
Lagrangian Function $L(x, \lambda, \nu) = f(x) + \sum_{i=1}^m \lambda_i g_i(x) + \sum_{i=1}^p \nu_i h_i(x)$							
Lagrange Dual Function (concave) $d(\lambda, \nu) = \inf_{x \in \text{dom}(f)} L(x, \lambda, \nu) \leq p^*$							
inf = infimum, the greatest lower bound of a set.							
<table border="0"> <tr> <th colspan="2">Primal and Dual Problem</th></tr> <tr> <td>$\min_x f(x)$</td><td></td></tr> <tr> <td>(P) : subj. to $g_i(x) \leq 0$ $h_i(x) = 0$</td><td>(D) : $\max_{\lambda, \nu} d(\lambda, \nu)$ subj. to $\lambda \geq 0$</td></tr> </table>		Primal and Dual Problem		$\min_x f(x)$		(P) : subj. to $g_i(x) \leq 0$ $h_i(x) = 0$	(D) : $\max_{\lambda, \nu} d(\lambda, \nu)$ subj. to $\lambda \geq 0$
Primal and Dual Problem							
$\min_x f(x)$							
(P) : subj. to $g_i(x) \leq 0$ $h_i(x) = 0$	(D) : $\max_{\lambda, \nu} d(\lambda, \nu)$ subj. to $\lambda \geq 0$						
• $d(\lambda, \nu)$ always concave • (D) convex even if (P) not • $d^* \leq p^* \rightsquigarrow d(\lambda, \nu)$ gens • Point (λ, ν) dual feas. if $\lambda \geq 0, (\lambda, \nu) \in \text{dom}(d)$ lower bound for p							
LP – Dual (P) : $\min_{x \in \mathbb{R}^n} c^T x$ subj. to $Ax = b, \quad Cx \leq e$ (D) : $\max_{\lambda, \nu} -b^T \nu - e^T \lambda, \quad \text{s.t. } A^T \nu + C^T \lambda + c = 0, \lambda \geq 0$							
QP – Dual with $Q \succ 0$ (P) : $\min_{x \in \mathbb{R}^n} \frac{1}{2} x^T Q x + c^T x, \quad \text{subj. to } Cx \leq e$ (D) : $\max_{\lambda, \nu} \frac{1}{2} \lambda^T C Q^{-1} C^T \lambda + (C Q^{-1} c + e)^T \lambda + \frac{1}{2} c^T Q^{-1} c$ subj. to $\lambda \geq 0$							
QP – Lagrangian <table border="0"> <tr> <td>$\min_x \frac{1}{2} x^T H x + q^T x + r$</td><td>$L = \frac{1}{2} x^T H x + q^T x + r$</td></tr> <tr> <td>s.t. $Gx \leq h$</td><td>$+ \lambda^T (Gx - h) + \nu^T (Ax - b)$</td></tr> <tr> <td>$Ax = b$</td><td>$+ \nabla_x L = Hx + q + G^T \lambda + A^T \nu$</td></tr> </table>		$\min_x \frac{1}{2} x^T H x + q^T x + r$	$L = \frac{1}{2} x^T H x + q^T x + r$	s.t. $Gx \leq h$	$+ \lambda^T (Gx - h) + \nu^T (Ax - b)$	$Ax = b$	$+ \nabla_x L = Hx + q + G^T \lambda + A^T \nu$
$\min_x \frac{1}{2} x^T H x + q^T x + r$	$L = \frac{1}{2} x^T H x + q^T x + r$						
s.t. $Gx \leq h$	$+ \lambda^T (Gx - h) + \nu^T (Ax - b)$						
$Ax = b$	$+ \nabla_x L = Hx + q + G^T \lambda + A^T \nu$						
Weak & Strong Duality Weak Duality – it is always true that $d^* \leq p^*$ Stront Duality – it is sometimes true that $d^* = p^*$ • Strong duality usually does not hold for non-convex problems • Can impose conditions on convex prob. to guarantee $d^* = p^*$ • Sometimes the dual much easier to solve than the primal • LP always has strong duality							

3.5.2	Slater Condition
If \exists at least one strictly feasible point $i \in$	
$\{x \mid Ax = b, g_i(x) < 0 \ \forall i\} \neq \emptyset \Rightarrow p^* = d^*$	
3.5.3	KKT Conditions
(1)	Primal Feasibility
$g_i(x^*) \leq 0, \ i = 1 \dots m \quad h_i(x^*) = 0, \ i = 1 \dots p$	
(2)	Dual feasibility
$\lambda^* \geq 0$	
(3)	Complementary Slackness
$\lambda_i^* g_i(x^*) = 0 \quad i = 1 \dots m$	
(4)	Stationarity
$\nabla L = \nabla f(x^*) + \sum_{i=1}^m \lambda_i^* \nabla g_i(x^*) + \sum_{i=1}^p \nu_i^* \nabla h_i(x^*) = 0$	

General Optimization Necessary condition
 x^*, λ^*, ν^* sol'n to $(P), (D)$ with 0 duality gap $\Rightarrow x^*, \lambda^*, \nu^*$ satisfy KKT
Convex Optimization Sufficient condition
 x^*, λ^*, ν^* satisfy KKT $\Rightarrow x^*, \lambda^*, \nu^*$ sol'n to $(P), (D)$ with 0 duality gap
Convex Opt. + Slater, Necessary & Sufficient condition
 If Slater's cond. holds, x^*, λ^*, ν^* are sol'n to $(P), (D)$ **IFF** KKT satisfied
Remark for convex opt. problem, KKT conditions sufficient \rightsquigarrow if x^*, λ^*, ν^* satisfy KKT then $p^* = d^*$

4	CFTOC
Constrained Finite Time Optimal Control	
$J^*(x(k)) = \min_U l_f(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i)$	
subj. to $x_{i+1} = Ax_i + Bu_i$	
$x_i \in \mathcal{X}, \quad u_i \in \mathcal{U}$	
$x_N \in \mathcal{X}_f, \quad x_0 = x(k)$	

- Quad. Cost / Squared Euclidian Norm:**

$$J(x(k)) = x_N^\top P x_N + \sum_{i=0}^{N-1} x_i^\top Q x_i + u_i R u_i$$

- p-Norm:** $J(x(k)) = ||Px_N||_p + \sum_{i=0}^{N-1} ||Qx_i||_p + ||Ru_i||_p$

4.1	Transform CFTOC to QP
QP Problem	
Goal – Rewrite Quad. Cost CF- TOC as QP	
\rightsquigarrow easier to solve	
4.1.1	Construction with Substitution, dense (good for large n)
Idea – Sub. state eqns $x_{i+1} = Ax_i + Bu_i, \quad x_0 = x(k)$	
Cost – Rewrite as (see Batch Approach for H and F)	

$$J^*(x(k)) = \min_U [U^\top \quad x(k)^\top] \begin{bmatrix} H & F \\ F^\top & Y \end{bmatrix} [U^\top \quad x(k)^\top]^\top$$

subj. to $GU \leq w + Ex(k)$

Constraints – Rewrite as $GU \leq w + Ex(k)$
 $\mathcal{X} = \{x | Ax x \leq b_x\} \quad \mathcal{U} = \{u | Au u \leq b_u\} \quad \mathcal{X}_f = \{x | A_f x \leq b_f\}$

$$G = \begin{bmatrix} A_u & \dots & A_u \\ 0 & \dots & 0 \\ A_x B & A_x A & A_x B \\ A_x A & A_x B & 0 \\ A_f A & A_f B & A_f B \end{bmatrix}, E = \begin{bmatrix} 0 \\ \dots \\ 0 \\ -A_x A \\ -A_x A \\ \dots \\ -A_f A \end{bmatrix}, w = \begin{bmatrix} b_u \\ \dots \\ b_u \\ b_x \\ b_x \\ \dots \\ b_f \end{bmatrix}$$

Solution For a given $x(k), U^*$ can be found via QP solver
4.1.2 Construction without Substitution, sparse (good for large N)
Idea – Keep state eqns as eq. constraints

$$\text{Cost with } z = \begin{bmatrix} x_1^\top & \dots & x_N^\top & u_0^\top & \dots & u_{N-1}^\top \end{bmatrix}^\top$$

$$J^*(x(k)) = \min_z \left[z^\top x(k)^\top \right] \begin{bmatrix} \bar{H} & 0 \\ 0 & Q \end{bmatrix} \begin{bmatrix} z^\top & x(k)^\top \end{bmatrix}^\top$$

subj. to $G_{in} z \leq w_{in} + E_{in} x(k)$

$$G_{eq} z = E_{eq} x(k)$$

$$\bar{H} = \text{diag}(Q, \dots, Q, P, R, \dots, R)$$

Equality Constraints from System Dyn. $x_{i+1} = Ax_i + Bu_i$

$$G_{eq} = \begin{bmatrix} \mathbb{I} & & \\ -A & \mathbb{I} & \\ \dots & A & \mathbb{I} \end{bmatrix} \begin{bmatrix} -B \\ -B \\ \dots \\ -B \end{bmatrix}, E_{eq} = \begin{bmatrix} A \\ 0 \\ \dots \\ 0 \end{bmatrix}$$

Inequality Constraints
 $\mathcal{X} = \{x | Ax x \leq b_x\} \quad \mathcal{U} = \{u | Au u \leq b_u\} \quad \mathcal{X}_f = \{x | A_f x \leq b_f\}$

$$G_{in} = \left[\begin{array}{c|ccc} 0 & & & \\ \hline A_x & A_x & \dots & \\ & & A_f & \\ \hline 0 & & & 0 \end{array} \right] \begin{array}{c|ccc} 0 & & & \\ \hline & 0 & \dots & \\ \hline A_u & A_u & \dots & A_u \end{array}, w_{in} = \begin{bmatrix} b_x \\ b_x \\ b_x \\ \dots \\ b_f \\ b_u \\ \dots \\ b_u \end{bmatrix},$$

$$E_{in} = \begin{bmatrix} -A_x^T & 0 & \dots & 0 \end{bmatrix}^\top$$

4.1.3	QP Feedback Solution
CFTOC problem as multiparametric QP	
$J^*(x(k)) = \min_U \left[U^\top x(k)^\top \right] \begin{bmatrix} H & F \\ F^\top & Y \end{bmatrix} \left[U^\top x(k)^\top \right]^\top$	
subj. to $GU \leq w + Ex(k)$	
Solution Properties	
● First component of optimal solution:	
$u_0^* = \kappa(x(k)), \quad \forall x(k) \in \mathcal{X}_0$	
$\kappa: \mathbb{R}^n \rightarrow \mathbb{R}^m$ is cont. and pw. affine on Polyhedra	
$\kappa(x) = F^j x + g^j \quad \text{if} \quad x \in CR^j, \quad j = 1, \dots, N^r$	
● Polyhedral sets $CR^j = \{x \in \mathbb{R}^n \mid H^j x \leq K^j\}, j = 1, \dots, N^r$	
are partition of the feasible polyhedron \mathcal{X}_0 .	
● Value func. $J^*(x(k))$ is convex and pw quad. on polyhedra.	

4.1.4	Transform p-norm CFTOC to LP
ℓ_∞ -Minimization	
$\min_{x \in \mathbb{R}^m} x _\infty \iff \min_{x, t} t$	
subj. to $Fx \leq g \quad \text{subj. to} \quad -1t \leq x \leq 1t, Fx \leq g$	
$-1t \leq x \leq 1t$ bounds abs value of every elem. with scalar t	
ℓ_1 -Minimization	
$\min_{x \in \mathbb{R}^m} x _1 \iff \min_{x \in \mathbb{R}^m, t \in \mathbb{R}^m} \mathbf{1}^\top t$	
subj. to $Fx \leq g \quad \text{subj. to} \quad -t \leq x \leq t, Fx \leq g$	
$ x _1 = \sum_{i=1}^n x_i \leq \sum_{i=1}^n t_i = \mathbf{1}^\top t \rightsquigarrow -t \leq x \leq t$	
bounds abs value of each component of x with a component of t	

4.1.5	Construction of ∞ -norm
Cost (with substitution)	
$\min_z \epsilon_N^x + \sum_{i=0}^{N-1} \epsilon_i^x + \epsilon_u^u$	
subj. to $-1_n \epsilon_i^x \leq \pm Q \left[A^i x_0 + \sum_{j=0}^{i-1} A^j Bu_{i-1-j} \right]$	
$-1_r \epsilon_N^x \leq \pm P \left[A^N x_0 + \sum_{j=0}^{N-1} A^j Bu_{N-1-j} \right]$	
$-1_m \epsilon_u^u \leq Ru_i$	
$x_i \in \mathcal{X}, \quad u_i \in \mathcal{U}, \quad x_f \in \mathcal{X}_f, \quad x_0 = x(k)$	

Substitution: $z := \{\epsilon_0^x \dots \epsilon_N^x, \epsilon_0^u \dots \epsilon_{N-1}^u, u_0^\top \dots u_{N-1}^\top\} \in \mathbb{R}^s$
 $s := (m + 1)N + N + 1$ results in:

$$\min_z c^\top z \quad \text{subj. to} \quad \bar{G} z \leq \bar{w} + \bar{S} x(k)$$

$$\bar{G} = \begin{bmatrix} G_e & G_u \end{bmatrix}, \quad \bar{S} = \begin{bmatrix} S_e \\ S \end{bmatrix}, \quad \bar{w} = \begin{bmatrix} w_e \\ w \end{bmatrix}$$

Solution for given $x(k), U^*$ can be obtained via LP solver
4.1.6 LP State Feedback Solution

Multiparam-LP	$\min_z c^\top z$ subj. to $\bar{G} z \leq \bar{w} + \bar{S} x(k)$
Properties	
-First component of sol'n has form: $u_0^* = \kappa(x(0)), \forall x(k) \in \mathcal{X}_0$	
$\kappa: \mathbb{R}^n \rightarrow \mathbb{R}^m$ is cont. & pw affine on Polyhedra	
$\kappa(x) = F^j x + g^j$ if $x \in CR^j, \quad j = 1, \dots, N^r$	
-Polyhedral sets $CR^j = \{x \in \mathbb{R}^n \mid H^j x \leq K^j\}$ are partition of \mathcal{X}_0	
-In case of multiple optimizers, a pw affine control law exists	
- $J^*(x(0))$ is convex, pw linear on polyhedra	
Quad vs 1/∞-norm cost	
Solution is either (n = # opt. var., FS = feas. set.)	
Quadratic Cost	Linear Cost
-unique & in interior of FS (no -Unbounded constraints active)	-unique at vertex of FS (at least n active cons-
-unique & on boundary of FS traints)	
(at least 1 const. active)	-multiple optima (min. 1 active const.)

5	MPC vs Classical Control										
5.1	Difference to Classical Control										
<table border="1"> <tr> <td>Classical Control main issues:</td><td>MPC main issues:</td></tr> <tr> <td>Disturbance rejections</td><td>Control constraints (input limits)</td></tr> <tr> <td>Noise insensitivity</td><td>Process/state constraints</td></tr> <tr> <td>Model uncertainty</td><td>(safety and physical constraints)</td></tr> <tr> <td>Usually in frequency domain</td><td>Usually in time domain</td></tr> </table>		Classical Control main issues:	MPC main issues:	Disturbance rejections	Control constraints (input limits)	Noise insensitivity	Process/state constraints	Model uncertainty	(safety and physical constraints)	Usually in frequency domain	Usually in time domain
Classical Control main issues:	MPC main issues:										
Disturbance rejections	Control constraints (input limits)										
Noise insensitivity	Process/state constraints										
Model uncertainty	(safety and physical constraints)										
Usually in frequency domain	Usually in time domain										

MPC can better handle constraints as they are implemented into the control scheme. Classical controllers usually us ad hoc constraint management or suboptimal operation.

5.2	Advantages & Challenges
Advantages:	
● Systematic and proper handling of constraints	
● High performance controller	
Challenges:	
● Implementation: \rightarrow real-time solving is challenging	
● Feasibility: Optimization problem may become infeasible in the future	
● Stability: Closed-loop stability is not automatically guaranteed	
● Robustness: Closed-loop system is not necessarily robust against uncertainties or disturbances	

6	Invariance
System	
Autonomous $x(k+1) = g(x(k))$	
Closed-Loop $x(k+1) = g(x(k), \kappa(x(k)))$ for given κ	
Positively Invariant Set (Minkowski sum of invariant sets is also invariant)	
Set \mathcal{O} positively invariant for autonomous system if $x(k) \in \mathcal{O} \Rightarrow x(k+1) \in \mathcal{O}, \quad \forall k \in \{0, 1, \dots\}$	
Maximal Positively Invariant Set	
$\mathcal{O}_\infty \subset \mathcal{X}$ positively invariant and contains all other \mathcal{O}	Initialize: $\Omega_0 \leftarrow \mathcal{X}$ Do: $\Omega_{i+1} \leftarrow \text{pre}(\Omega_i) \cap \Omega_i$ Until: $\Omega_{i+1} = \Omega_i \rightarrow \mathcal{O}_\infty = \Omega_i$
Pre-Set	
Given set S , the pre-set of S is the set of states that evolve into S in 1 k:	
$x(k+1) = g(x(k))$ $\Rightarrow \text{pre}(S) := \{x \mid g(x) \in S\}$	$x(k+1) = Ax(k)$ $\Rightarrow \text{pre}(S) := \{x \mid Ax \in S\}$

Pre-Set Computation Linear System
 Set $S := \{x \mid Fx \leq f\}, x(k+1) = Ax(k)$ then
 $\text{pre}(S) := \{x \mid Ax \in S\} = \{x \mid FAx \leq f\}$

- For $\{x \mid Fx \leq f\}$, if $F \downarrow$ or $f \uparrow \rightsquigarrow$ **Less Restrictive**
- $S \cap F \rightsquigarrow$ constraints from both sets active

Invariant Set Conditions
Given set S , the pre-set of S is the set of states that evolve into S in one time step. Set \mathcal{O} is positively invariant set iff $\mathcal{O} \subseteq \text{pre}(\mathcal{O}) \Leftrightarrow \text{pre}(\mathcal{O}) \cap \mathcal{O} = \mathcal{O}$
Necessary if $\mathcal{O} \not\subseteq \text{pre}(\mathcal{O})$, then $\exists \bar{x} \in \mathcal{O}$ s.t $\bar{x} \notin \text{pre}(\mathcal{O}) \rightsquigarrow \bar{x} \in \mathcal{O}, g(\bar{x}) \notin (\mathcal{O})$, thus \mathcal{O} not positively invariant
Sufficient if \mathcal{O} not pos invar set, then $\exists \bar{x} \in \mathcal{O}$ s.t $g(\bar{x}) \notin \mathcal{O} \rightsquigarrow \bar{x} \in \mathcal{O}, \bar{x} \notin \text{pre}(\mathcal{O})$ thus $\mathcal{O} \not\subseteq \text{pre}(\mathcal{O})$

6.1	Control Invariance
Control Invariant Set	
Set $C \subseteq \mathcal{X}$ control invariant if $x(k) \in C \Rightarrow \exists u(k) \in \mathcal{U}$ s.t $g(x(k), u(k)) \in C \ \forall k \in \mathbb{N}^+$	

Maximal Control Invariant Set
Set C_∞ maximal control invariant if it is control invariant and contains all control invariant sets contained in \mathcal{X}

For all states in C_∞ , there exists control law s.t system constraints never violated \rightsquigarrow **The best any controller could ever do**
Pre-Set $\text{pre}(S) := \{x \mid \exists u \in \mathcal{U}$ s.t $g(x, u) \in S\}$

Control Invariant Set \Rightarrow Control Law
Let C be the control invariant set for $x(k+1) = g(x(k), u(k))$ The control law $\kappa(x(k))$ will guarantee that the system satisfies constraints $\forall t$ if $g(x, \kappa(x)) \in C \ \forall x \in C \rightsquigarrow$ With f as any function, synthesize control law $\kappa: \kappa(x) := \text{argmin}\{f(x, u) \mid g(x, u) \in C\}$

- Does not ensure sys. will converge, but will satisfy constraints
- Don't often do because calculating control invariant sets is very hard
- MPC implicitly** describes cont. invar. set s.t easy to represent/compute

6.2	Practical Invariant Set Computation
Minkowski-Weyl Theorem	
For $P \subset \mathbb{R}^d$ following statements equivalent:	
● P polytope, $\exists A, b$ s.t $P = \{x \mid Ax \leq b\}$	
● P finitely generated, \exists finite set of $\{v_i\}$ s.t $P = \text{co}(\{v_1 \dots v_s\})$	

6.2.1	Invariant Sets from Lyapunov Functions
Lemma If $V: \mathbb{R}^n \rightarrow \mathbb{R}$ a Lyap. func. for sys. $x(k+1) = g(x(k))$, then $Y := \{x \mid V(x) \leq \alpha\}$ is an invariant set for all $\alpha \geq 0$	
Proof: $V(g(x)) - V(x) < 0 \rightsquigarrow$ once $V(x(k)) \leq \alpha$, stays there $\forall j \geq k$	
6.2.2	Maximum Ellipsoidal Invariant Sets

For $x(k+1) = Ax(k)$ with $P \succ 0$ with $A^\top PA - P \prec 0$ then $V(x(k)) = x(k)^\top Px(k)$ is Lyap. function. Find largest α s.t set $Y_\alpha \in \mathcal{X}$
 $Y_\alpha := \{x \mid x^\top Px \leq \alpha\} \subset \mathcal{X} := \{x \mid Fx \leq f\}$
 Equivalent to $\max_\alpha \alpha$ subj. to $h_{Y_\alpha}(F_i) \leq f_i \ \forall i \in \{1 \dots n\}$

Support of an ellipse: $h_{Y_\alpha}(F_i) = \max_x F_i x$ subj. to $x^\top Px \leq \alpha$
 F_i and f_i are the rows of the polytopic description of \mathcal{X} and \mathcal{U}

Change of Variables: $y := P^{1/2} x$
 $\rightsquigarrow h_{Y_\alpha}(F_i) = \max_x F_i P^{-1/2} y \quad \text{s.t} \quad y^\top y \leq \sqrt{\alpha}^2$

Maximizer found by inspection:

$$h_{Y_\alpha}(F_i) = F_i P^{-1/2} \frac{P^{-1/2} F_i^\top}{||P^{-1/2} F_i^\top||} \sqrt{\alpha} = ||P^{-1/2} F_i^\top|| \sqrt{\alpha}$$

Largest ellipse now 1-dim optimization problem:

$$\alpha^* = \max_\alpha \alpha \quad \text{s.t.} \quad ||P^{-1/2} F_i^\top||^2 \alpha \leq f_i^2 \ \forall i \in \{1 \dots n\}$$

$$= \min_{i \in \{1 \dots n\}} \frac{f_i^2}{F_i P^{-1} F_i^\top}$$

7	MPC Formulation
System: $x(k+1) = Ax(k) + Bu(k), x(k) \in \mathbb{R}^n, u(k) \in \mathbb{R}^m$	
Control Law is defined by: $u = u^*(0)$	
$J^*(x(k)) = \min_U l_f(x_N) + \sum_{i=1}^{N-1} l(x_i, u_i)$	
subj. to $x_{i+1} = Ax_i + Bu_i$	
$x_i \in \mathcal{X}, \quad u_i \in \mathcal{U}$	
$x_N \in \mathcal{X}_f, \quad x_0 = x(k)$	

Assumptions that need to be met:

- Stage cost pos def, strictly positive, only 0 at origin
- (a) Terminal set invariant under local control law $\kappa_f(X)$:
 $x_{i+1} = Ax_i + B\kappa_f(x_i)$
(b) All state and input constraints satisfied in \mathcal{X}_f
- Terminal cost is cont. Lyap. func. in terminal set \mathcal{X}_f and satisfies
 $l_f(x_{i+1}) - l_f(x_i) \leq -l(x_i, \kappa_f(x_i)) \quad \forall x_i \in \mathcal{X}_f$

If 1-3 are met: CL system under MPC control law $u_0^*(x)$ asympt. stable and set \mathcal{X}_N is positive invariant for system $x(k+1) = Ax(k) + Bu_0^*(x(k))$
Often Quadratic Cost:

$$J^*(x(k)) = \min x_N^\top P x_N + \sum_{i=0}^{N-1} x_i^\top Q x_i + u_i^\top R u_i$$

subj. to $x_{i+1} = Ax_i + Bu_i$

$$x_i \in \mathcal{X}, \quad u_i \in \mathcal{U}, \quad x_N \in \mathcal{X}_f, \quad x_0 = x(k)$$

$$Q = Q^\top \succeq 0, \quad R = R^\top \succ 0$$

For 3. this implies: $A_{cl} = A + BK$

$$A_{cl}^\top P A_{cl} - P \preceq -Q(-K^\top R K), \quad A_{cl}^\top P B_{cl} - P \preceq 0$$

7.1	Loss Of Feasibility & Stability
Infinite-Horizon Solve RHC for $N = \infty$, OL traj. are same as CL traj.	
● If problem feasible, CL trajectories always feasible	
● If cost finite, states and inputs will converge asympt. to origin	
Finite-Horizon RHC "short-sighted" approximating ∞ -horizon controller	
● Feasibility – after some steps finite horizon optimal control problem may become infeasible (disturbances, model mismatch)	
● Stability – generated inputs may not lead to traj. that converge to origin	
Solution Introduce terminal cost & constraints to ensure feas. & stab.	
7.2	Feasibility & Stability Guarantees

Proof Strategy
Recursive Feasibility show existence of feasible control sequence for all time when starting from feasible initial point

- Assume feas. of $x(k), \{u_0^*, \dots, u_{N-1}^*\}, \{x_0^*, \dots, x_N^*\}$
- At $x(k+1) \Rightarrow \{u_1^*, \dots, \kappa_f(x_N^*)\}$ should be feas.

Stability show that optimal cost is Lyapunov function

- l_f necessary to provide cost decrease for asympt. stability

Terminal Constraint At Zero $x_N \in \mathcal{X}_f = 0$
If at 0 and no input is given system stays there \rightsquigarrow stable and feasibly point. need large N to approximate maximum control invariant set
General Terminal Set \mathcal{X}_f
Need assumptions 1-3 for stability guarantees. Cost decrease proof:

$$J^*(x(k+1)) \leq \sum_{i=1}^{N-1} l(x_i^*, u_i^*) + \underbrace{l(x_N^*, u_N^* = \kappa_f(x_N^*))}_{\text{stage cost at } k+1} + \underbrace{l_f(Ax_N^* + B\kappa_f(x_N^*))}_{\text{cost of propagated state}}$$

$$= \underbrace{\sum_{i=0}^{N-1} l(x_i^*, u_i^*)}_{J^*(x(k)) - l_f(x_N^*)} - l(x_0^*, u_0^*) + l(x_N^*, \kappa_f(x_N^*)) + l_f(Ax_N^* + B\kappa_f(x_N^*))$$

$$= J^*(x(k)) - l(x(k), u_0^*) + l(x_N^*, \kappa_f(x_N^*)) + l_f(Ax_N^* + B\kappa_f(x_N^*)) - l_f(x_N^*)$$

$\underbrace{\hspace{10em}}_{\text{subtract cost at stage } k > 0} \leq 0$ by assumption of Lyap func.in terminal set \mathcal{X}_f

8 Practical MPC
8.1 MPC Reference Tracking
8.1.1 Steady-State Target Tracking
Target Condition
$\begin{matrix} x_s = Ax_s + Bu_s \\ z_s = Hx_s = r \end{matrix} \iff \underbrace{\begin{bmatrix} \mathbb{I} - A & -B \\ H & 0 \end{bmatrix}}_{(n_x+n_r) \times (n_x+n_u)} \begin{bmatrix} x_s \\ u_s \end{bmatrix} = \begin{bmatrix} 0 \\ r \end{bmatrix}$
<ul style="list-style-type: none"> In general, (x_s, u_s) must satisfy them In case of multiple feas. u_s, compute 'cheapest' <ul style="list-style-type: none"> $\min u^T R_s u_s$, subj. to [Target Condition], $x_s \in \mathcal{X}, u_s \in \mathcal{U}$ In general, assume target problem is feasible If no sol'n \exists: compute reachable point 'closest' to r <ul style="list-style-type: none"> $\min(Hx_s - r)^T Q_s (Hx_s - r), \quad \text{subj. to } x_s = Ax_s + Bu_s$
8.1.2 Reference Tracking
MPC Design
$\min_U z_N - Hx_s _{P_z}^2 + \sum_{i=1}^{N-1} z_i - Hx_s _{Q_z}^2 + u_i - u_s _{R}^2$
subj. to [model, constraints], $x_0 = x(k)$
Delta Formulation
Set pt. tracking $\xrightarrow{\text{Coord.Trans.}}$ Regulation Problem
$\begin{matrix} \Delta x := x - x_s & \Big & G_x \Delta x \leq h_x - G_x x_s \\ \Delta u := u - u_s & & G_u \Delta u \leq h_u - G_u u_s \end{matrix}$
<ul style="list-style-type: none"> Obtain target steady-state corresponding to reference r Initial state $\Delta x(k) = x(k) - x_s$ Apply reg problem to new system in Δ-Formulation <ul style="list-style-type: none"> $\min \left[V_f(\Delta x_N) + \sum_{i=1}^{N-1} \Delta x_i^T Q \Delta x_i + \Delta u_i^T R \Delta u_i \right]$ subj. to $\Delta x_{i+1} = A \Delta x_i + B \Delta u_i, \quad G_x \Delta x_i \leq h_x - G_x x_s$ $G_u \Delta u_i \leq h_u - G_u u_s, \quad \Delta x_N \in \mathcal{X}_f, \quad \Delta x_0 = \Delta x(k)$ Find optimal sequence of ΔU^* Input applied to system $u_0^* = \Delta u_0^* + u_s$
Convergence
Assume target feasible with $x_s \in \mathcal{X}, u_s \in \mathcal{U}$, choose terminal weight $V_f(x)$ and constraint \mathcal{U}_f as in regulation case satisfying
<ul style="list-style-type: none"> $\mathcal{X}_f \subseteq \mathcal{X}, Kx \in \mathcal{U} \quad \forall x \in \mathcal{X}_f$ $V_f(x(k+1)) - V_f(x(k)) \leq -l(x(k), Kx(k)) \quad \forall x \in \mathcal{X}_f$
If in addition the target reference x_s, u_s is such that
<ul style="list-style-type: none">$x_s \oplus \mathcal{X}_f \subseteq \mathcal{X}, \quad K\Delta x + u_s \in \mathcal{U}, \quad \forall \Delta x \in \mathcal{X}_f$
then CL system converges to target reference
$x(k) \rightarrow x_s, z(k) = Hx(k) \xrightarrow{k \rightarrow \infty} r$
Proof
<ul style="list-style-type: none"> Invariance under local ctrl law inherited from regulation case Constraint satisfaction provided by extra conditions <ul style="list-style-type: none"> $-x_s \oplus \mathcal{X}_f \subseteq \mathcal{X} \rightarrow x \in \mathcal{X} \forall \Delta x \in \mathcal{X}_f$ $-K\Delta x + u_s \in \mathcal{U} \forall \Delta x \in \mathcal{X}_f \rightarrow u \in \mathcal{U}$ Fron asympt stability of the regulation problem: $\Delta x(k) \xrightarrow{k \rightarrow \infty} 0$
Terminal Set
<ul style="list-style-type: none"> Set of feasible targets may be significantly reduced. Enlarge set of feasible targets by scaling terminal set for regulation $\mathcal{X}_f^{\text{scaled}} = \alpha \mathcal{X}_f$ Invariance maintained if \mathcal{X}_f invariant \rightsquigarrow so is $\alpha \mathcal{X}_f$ Choose α s.t. x, u constraints still satisfied \rightsquigarrow scaling target dependent Targets at the boundary of the constraints: $x_N = x_s$, corresponds to 0-terminal set in regulation case
8.2 Disturbance Rejection
Augmented Model
$\begin{matrix} x_{k+1} = Ax_k + Bu_k + B_d d_k \\ d_{k+1} = d_k, \quad y_k = Cx_k + C_d d_k \end{matrix}$
Observability of aug. system: $\text{rank} \left(\begin{bmatrix} A - \mathbb{I} & B_d \\ C & C_d \end{bmatrix} \right) \stackrel{!}{=} n_x + n_d$
Inuition At steady-state $\begin{bmatrix} A - \mathbb{I} & B_d \\ C & C_d \end{bmatrix} \begin{bmatrix} x_s \\ d_s \end{bmatrix} = \begin{bmatrix} 0 \\ y_s \end{bmatrix}$ and given y_s, d_s must be uniquely defined
Linear State Estimation
Observer For Augmented Model:
$\begin{bmatrix} \hat{x}(k+1) \\ \hat{d}(k+1) \end{bmatrix} = \begin{bmatrix} A & B_d \\ C & C_d \end{bmatrix} \begin{bmatrix} \hat{x}(k) \\ \hat{d}(k) \end{bmatrix} + \begin{bmatrix} B \\ 0 \end{bmatrix} u(k)$
$+ \begin{bmatrix} L_x \\ L_d \end{bmatrix} (-y(k) + C\hat{x}(k) + C_d \hat{d}(k))$
Error Dynamics \rightsquigarrow choose L s.t error dynamics asympt. stable
$\begin{bmatrix} x(k+1) - \hat{x}(k+1) \\ d(k+1) - \hat{d}(k+1) \end{bmatrix} = \begin{bmatrix} A & B_d \\ 0 & \mathbb{I} \end{bmatrix} \begin{bmatrix} x(k) - \hat{x}(k) \\ d(k) - \hat{d}(k) \end{bmatrix}$
$- \begin{bmatrix} L_x \\ L_d \end{bmatrix} (C\hat{x}(k) + C_d \hat{d}(k) - Cx(k) - C_d d(k))$
$= \left(\begin{bmatrix} A & B_d \\ 0 & \mathbb{I} \end{bmatrix} + \begin{bmatrix} L_x \\ L_d \end{bmatrix} [C \quad C_d] \right) \begin{bmatrix} x(k) - \hat{x}(k) \\ d(k) - \hat{d}(k) \end{bmatrix}$

Observer State-Space
Suppose observer asympt. stable and $n_y = n_d$
$\begin{bmatrix} A - \mathbb{I} & B \\ C & 0 \end{bmatrix} \begin{bmatrix} \hat{x}_\infty \\ u_\infty \end{bmatrix} = \begin{bmatrix} -B_d \hat{d}_\infty \\ y_\infty - C_d \hat{d}_\infty \end{bmatrix}$
\rightsquigarrow Observer output $C\hat{x}_\infty + C_d \hat{d}_\infty$ tracks y_∞ without offset
Offset-Free Tracking
Goal Track constant r : $z(k) = Hy(k) \rightarrow r$ as $k \rightarrow \infty$
Steady-State Condition
$x_s = Ax_s + Bu_s + B_d \hat{d}_\infty, \quad z_s = H(Cx_s + C_d \hat{d}_\infty) = r$
<ul style="list-style-type: none"> Best forecast for d_∞ is current estimate $\hat{d}_\infty = \hat{d}$ Same Procedure for regulation case with $r = 0$
Offset-Free Tracking Condition: $\begin{bmatrix} A - \mathbb{I} & B \\ HC & 0 \end{bmatrix} \begin{bmatrix} x_s \\ u_s \end{bmatrix} = \begin{bmatrix} -B_d \hat{d} \\ r - HC_d \hat{d} \end{bmatrix}$
Offset-Free Tracking Procedure
<ol style="list-style-type: none"> Estimate \hat{x} & \hat{d} Obtain (x_s, u_s) from steady-state tgt problem using \hat{d} Solve MPC problem for tracking using $\hat{d}, \hat{x}_i := x_i - x_s, \hat{u}_i = u_i - u_s$ $\min_U V_f(\tilde{x}_N) + \sum_{i=0}^{N-1} (\tilde{x}_i)^T Q(\tilde{x}_i) + (\tilde{u}_i)^T R(\tilde{u}_i)$
subj. to $x_{i+1} = Ax_i + Bu_i + B_d \hat{d}_i, \quad d_{i+1} = \hat{d}_i$
$x_i \in \mathcal{X}, \quad u_i \in \mathcal{U}, \quad x_0 = \hat{x}(k), \quad d_0 = \hat{d}(k), \quad x_n - x_s \in \mathcal{X}_f$
Offset-Free Tracking: Main Result
With $u_0^* = \kappa(\hat{x}(k), \hat{d}(k), r)$. Assuming $n_d = n_y$, RHC recursively feasible and unconstrained for $k \geq j, j \in \mathbb{N}^+$ and the CL system:
$x(k+1) = Ax(k) + B\kappa(\hat{x}(k), \hat{d}(k), r) + B_d \hat{d}$
$\hat{x}(k+1) = (A + L_x C)\hat{x}(k) + (B_d + L_x C_d)\hat{d}(k) + B\kappa(\hat{x}(k), \hat{d}(k), r) - L_x y(k)$
$\hat{d}(k+1) = L_d C\hat{x}(k) + (\mathbb{I} + L_d C_d)\hat{d}(k) - L_d y(k)$
converging, i.e. $((\hat{x}, \hat{d}) \xrightarrow{k \rightarrow \infty} (x_\infty, d_\infty))$
Then $z(k) = Hy(k) \xrightarrow{k \rightarrow \infty} r$
8.3 Enlarging Feasible Set
8.3.1 No Terminal Set
Motivation Terminal constraints reduce feasible set, Stability guarantees can add large number of constraints and adds state constraints to problems with only input constraints.
Goal MPC without terminal constraints with guaranteed stability
Note Feasible set without terminal constraint not invariant
MPC Without Terminal Set
Can remove terminal constraint while maintaining stability if
<ul style="list-style-type: none"> Initial state lies in sufficiently small subset of feasible set N sufficiently large
s.t term. state satisfies term. const. without envorncing it in the optimization.
\rightsquigarrow Sol'n of finite-horizon MPC problem corresponds to ∞ -horizon sol'n
Advantage – Controller defined in larger feasible set
Disadvantage – Characterization of region of attaction of specification of required horizon length extremely difficult
<ul style="list-style-type: none"> Term constraint provides sufficient condition for stab: Region of attraction without term constraint may be larger than with In practice: Enlarge horizon and check stability by sampling $N \uparrow \rightsquigarrow$ RoA approaches max control invar. set
8.3.2 Soft constraints
Motivation Input constraints usually 'hard' due to physical limits, state constraints rarely 'hard' (more safety and comfort reasons)
Goal Min size & duration of violation (usually conflict!)
MPC Problem Setup
$\min_{u_N} x_N^T P x_N + l_\epsilon(\epsilon_N) + \sum_{i=0}^{N-1} x_i^T Q x_i + u_i^T R u_i + l_\epsilon(\epsilon_i)$
s.t. $x_i = Ax_i + Bu_i, Hx_i \leq k_x + \epsilon_i, H_u u_i \leq k_u, \quad \epsilon_i \geq 0$
Requirement on l_ϵ
Original Problem $\min_z f(z) \quad \text{s.t. } g(z) \leq 0$
“Softened” Problem $\min_z f(z) + l_\epsilon(\epsilon) \quad \text{s.t. } g(z) \leq \epsilon, \epsilon \geq 0$
If original problem has feasible solution z^* , Softened problem should have same solution z^* , and $\epsilon = 0$.
Note $l_\epsilon(\epsilon_i) = s\epsilon_i^2$ does not fulfill requirement
Choice of Penalty
<ul style="list-style-type: none"> Quad. Penalty $l_\epsilon(\epsilon_i) = \epsilon_i^T S \epsilon_i$ (e.g. $S = Q$) Quad. + Linear Penalty $l_\epsilon(\epsilon_i) = \epsilon_i^T S \epsilon_i + v \epsilon_i _1 / \infty$
Exact Penalty Function
$l_\epsilon(\epsilon) = v \cdot \epsilon$ satisfies requirement for any $v > \lambda^* \geq 0$, where λ^* is optimal Lagrange multiplier for original problem.
In practice combined cost \rightarrow exact penalty and tuning capabilities
$l_\epsilon(\epsilon) = v \cdot \epsilon + \epsilon^T S \epsilon$
with $v > \lambda^*$ and $S \succ 0$.
Tuning
<ul style="list-style-type: none"> Increasing S leads to hardeing of constraints \rightarrow reduced violation size but longer duration Increasing v leads to constraint satisfaction if possible \rightarrow larger but shorter violation

Objective Separation
1. Minimize violation over horizon:
$\epsilon^{\min} = \text{argmin}_{u, \epsilon} \sum_{i=0}^{N-1} \epsilon_i^T S \epsilon_i + v^T \epsilon_i$
s.t. $x_{i+1} = Ax_i + Bu_i Hx_i \leq k_x + \epsilon_i$
$H_u u_i \leq k_u, \quad \epsilon_i \geq 0$
2. Optimize Controller performance
$\min_{u_N} x_N^T P x_N + \sum_{i=0}^{N-1} x_i^T Q x_i + u_i^T R u_i$
s.t. $x_{i+1} = Ax_i + Bu_i, \quad Hx_i \leq k_x + \epsilon_i^{\min}, \quad H_u u_i \leq k_u$
Simplifies tuning and constraint satisfied if possible, but two optimization problems have to be solved.
Note SC MPC does not provide stability guarantee for OL unstable sys.
9 Robust MPC for Linear Systems
9.1 Robust Open-Loop MPC
9.1.1 Uncertainty Models
Motivation: Random noise w influences system evolution, Model structure is unknown, Unknown parameters θ impact dynamics.
Uncertain Constrained System
$x(k+1) = g(x(k), u(k), w(k); \theta),$
$x, u, w, \theta \in \mathcal{X}, \mathcal{U}, \mathcal{W}, \Theta$
9.1.2 Robust Invariance
Robust Positive Invariant Set
Set \mathcal{O}^W said to be robust positive invariant for the autonomous system $x(k+1) = g(x(k), w(k))$ if
$x \in \mathcal{O}^W \Rightarrow g(x, w) \in \mathcal{O}^W, \quad \forall w \in \mathcal{W}, \forall k$
Robust Pre Set
Given set Ω and dynamic system $x(k+1) = g(x(k), w(k))$,
$\text{pre}^W(\Omega) := \{x \mid g(x, w) \in \Omega \quad \forall w \in \mathcal{W}\}$
Maximal Robust Positively Invariant Set
$\mathcal{O}_\infty^W \subset \mathcal{X}$ positively invariant and contains all other \mathcal{O}^W .
Calculation using the algorithm for the nominal case.
Computing Robust Pre-Sets for Linear Systems
System $Ax(k) + w(k)$, set $\Omega := \{x \mid Fx \leq f\}$
$\text{pre}^W(\Omega) = \{x \mid FAx \leq f - \max_{w \in W} Fw\}$
$= \{x \mid FAx \leq f - h_{\mathcal{W}^F}(F)\}$
$h_{\mathcal{W}}$ is the support function
Robust Invariant Set Conditions
Set \mathcal{O}^W is a robust positive invariant set iff
$\mathcal{O}^W \subseteq \text{pre}^W(\mathcal{O}^W) \Leftrightarrow \text{pre}^W(\mathcal{O}^W) \cap \mathcal{O}^W = \mathcal{O}^W$
9.1.3 Impact of Additive Bounded Noise
Additive Bounded Noise System:
$x(k+1) = Ax(k) + Bu(k) + w(k),$
$x, u, w \in \mathcal{X}, \mathcal{U}, \mathcal{W}$
Uncertain State Evolution:
$\phi_i = A^i x_0 + \underbrace{\sum_{j=0}^{i-1} A^j Bu_{i-1-j}}_{x_i \equiv \text{Nominal System}} + \underbrace{\sum_{j=0}^{i-1} A^j w_{i-1-j}}_{\text{Disturbance Offset}}$
Robust Open-Loop MPC
Robust Open-Loop MPC
$\min_U l_f(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i)$
subj. to $x_{i+1} = Ax_i + Bu_i$
$x_i \in \mathcal{X} \ominus \left(\bigoplus_{j=0}^{i-1} A^j \mathcal{W} \right), \quad u_i \in \mathcal{U}$
$x_0 = x(k), \quad x_N \in \mathcal{X}_f \ominus \left(\bigoplus_{j=0}^{N-1} A^j \mathcal{W} \right)$
where $\mathcal{X}_f \subseteq \mathcal{X}$ robust positive invariant set for system $(A+BK)x(k) + w(k)$ with $w \in \mathcal{W} \quad \forall k$ for some stabilizing K , and $Kx \in \mathcal{U} \quad \forall x \in \mathcal{X}_f$
Intuition Nominal MPC, but with tighter state constraints
Open-Loop: Not accounting for FB during solving, just plan ahead for w
Caution: Unstable systems $A^{i-1} \mathcal{W}$ grows \rightarrow use 'pre-stabilization' $u_i = Kx_i + u_i$
Potentially very small region of attraction, particularly for unstable sys
9.1.4 Robust Constrained Control
Goals: Design $u(k) = \kappa(x(k))$ such that the system
<ol style="list-style-type: none"> Satisfies constraints: $\{x(k)\} \subset \mathcal{X}, \{u(k)\} \subset \mathcal{U}$ for all disturbances Is Stable: converges to a neighborhood of the origin Optimizes (expected/worst-case) 'Performance' Maximizes Set $\mathcal{X}(0)$ Condition 1-3 met

(a) Robust Constraint Satisfaction
Ensure all states $\phi_i(x_0, U, W)$ satisfy system constraints \mathcal{X} :
<ul style="list-style-type: none"> State & Input Constraints for $i = 0, \dots, N-1$, Enforce constraints explicitly by imposing: <ul style="list-style-type: none"> $\phi_i \in \mathcal{X}, u_i \in \mathcal{U}, \forall W \in \mathcal{W}^N$ Terminal Constraints for $i = N, \dots$. Enforce constraints implicitly by: <ul style="list-style-type: none"> Constraining $\phi_N \in \mathcal{R}$ robust invariant set \mathcal{X}_f and $K\mathcal{X}_f \in \mathcal{U}$ for $\phi_{i+1} = (A+BK)\phi_i + w_i$
We want for all $i = 0, \dots, N$:
$\phi_i(x_0, U, W) = \left\{ x_i + \sum_{j=0}^{i-1} A^j w_{i-1-j} \mid W \in \mathcal{W}^i \right\} \subseteq \mathcal{X}$
Assume $\mathcal{X} = \{x \mid Fx \leq f\}$ (polyhedron)
$Fx_i \leq f - h_{\mathcal{W}^F} \left(F \sum_{j=0}^{i-1} A^j \right)$
\rightarrow tightening constrains on the nominal system.
Support function $h_{\mathcal{W}^F}$ can be pre-computed offline.
Same goes for $i = N, \dots, \infty$, i.e. $\phi_N(x_0, U, W) \subseteq \mathcal{X}_f$.
Requirement can be rewritten as:
$\phi_i \in x_i \oplus (\mathcal{W} \oplus A\mathcal{W} \dots A^{i-1}\mathcal{W}) \subseteq \mathcal{X}$
or
$x_i \in \mathcal{X} \ominus \left(\bigoplus_{j=0}^{i-1} A^j \mathcal{W} \right)$
$\mathcal{F}_i = \bigoplus_{j=0}^{i-1} A^j \mathcal{W}$ is called disturbance reachable set.
Note: $\mathcal{F}_{i+1} = A\mathcal{F}_i \oplus \mathcal{W}$
Caution: Must ensure term state contained in robust invariant set
Intuition: Tightening constraints on the nominal system
(b) Is Stable: To show stability more general stability theory is needed.
(c) – Optimizes Performance
Cost to Minimize:
$J(x_0, U, W) := l_f(\phi_N(x_0, U, W)) + \sum_{i=0}^{N-1} l(\phi_i(x_0, U, W), u_i)$
Several options to eliminate dependence on W :
<ul style="list-style-type: none"> Minimize expected value: $J_N(x_0, U) = \mathbb{E}\{J(x_0, U, W)\}$ Take the worst case: $J_N(x_0, U) := \max_{W \in \mathcal{W}, N=1} J(x_0, U, W)$
<ul style="list-style-type: none"> Take the Nominal Case $J_N(x_0, U) := J(x_0, U, 0)$
(d) Maximizes Set: potentially very small region of attraction
9.2 Robust Closed Loop MPC
Increase the feasibly set using closed-loop feedback.
9.2.1 Closed-Loop Predictions
Goal optimize over seq. of funces $\{u_0, \mu_1(\cdot), \dots, \mu_{N-1}(\cdot)\}$
where $\mu_i(x_i) : \mathbb{R}^n \rightarrow \mathbb{R}^m$ is called control policy
Problem Can't optimize over arbitrary functions!
Solution assume some structure on functions μ_i
Pre-Stabilization $\mu_i(x_i) = Kx_i + v_i$
Fixed K , s.t $A+BK$ stable \rightarrow Simple, often conservative
Linear Feedback $\mu_i(x_i) = K_i x_i + v_i$
Optimize over K_i, v_i , Non-Convex – Extremely difficult to solve
Disturbance Feedback $\mu_i(x_i) = \sum_{j=0}^{i-1} M_{ij} w_j + v_i$
Optimize over $M_{ij}, v_i \rightsquigarrow$ Equiv to linear feedback but Convex \rightarrow Ef-fective, but computationally intense
Tube-MPC $\mu_i(x_i) = v_i + K(x_i - \bar{x}_i)$
Fixed K , s.t $A+BK$ stable \rightarrow Optimize over $\bar{x}_i, v_i \rightarrow$ Simple, can be effective
9.2.2 Tube-MPC
System: $x(k+1) = Ax(k) + Bu(k) + w(k) \quad x, u \in \mathcal{X}, \mathcal{U} \quad w \in \mathcal{W}$
Idea Separate available control authority into 2 parts
(1) Portion that steers nominal system to origin:
$z(k+1) = Az(k) + Bv(k)$
(2) Portion that compensates for deviations from this system $u_i = K(x_i - z_i) + v_i$ (keeps real traj close to nominal), for some linear K , which stabilizes nominal system
\rightarrow Fix linear feedback K offline and optimize over nominal trajectory $\{v_0, \dots, v_{N-1}\} \rightarrow$ convex problem
Error Dynamics
Define $e_i := x_i - z_i \rightsquigarrow e_{i+1} = (A+BK)e_i + w_i$
Bound maximum error, how far 'real' traj from nominal
$e_{i+1} = (A+BK)e_i + w_i \quad w_i \in \mathcal{W}$
Dynamics $A+BK$ are stable, set \mathcal{W} bounded \rightsquigarrow Set \mathcal{E} s.t e stays inside $\forall k \rightarrow$ 'minimal robust invariant set'
Tube-MPC Procedure
<ol style="list-style-type: none"> Compute set \mathcal{E} that error remains inside Modify constraints on nominal traj $\{z_i\}$ Formulate as convex optimization problem
(a) Minimum Robust Invariant Set (mRPI)
Minimum Robust Invariant Set
$F_\infty = \bigoplus_{j=0}^{\infty} A^j \mathcal{W}$
If $F_n = F_{n+1} \Rightarrow F_n = F_\infty$
Algorithm to Compute F_∞
<pre> $\Omega_0 \leftarrow \{0\}$ loop $\Omega_{i+1} \leftarrow \Omega_i \oplus A^i \mathcal{W}$ if $\Omega_{i+1} = \Omega_i$ then return $F_\infty = \Omega_i$ end if end loop </pre>
-Finite n does not always exist, 'large' n often good approximation.
-If not finite, other methods for small invariant sets, bit larger than F_∞

(b) **Global Nominal Trajectory Constraints**

Noisy System Trajectory:
Given nominal trajectory z_i noisy system trajectory $x_i = z_i + e_i$ will be somewhere in \mathcal{E}

$$x_i \in z_i \oplus \mathcal{E} = \{z_i + e \mid e \in \mathcal{E}\}$$

Goal $x_i, u_i \in \mathcal{X}, \mathcal{U}$ for all $\{w_i\} \in \mathcal{W}^J$

State Condition Necessary & Sufficient Condition

$$z_i \oplus \mathcal{E} \subseteq \mathcal{X} \Leftrightarrow z_i \in \mathcal{X} \oplus \mathcal{E}$$

Input Condition:

$$u_i \in K\mathcal{E} \oplus v_i \subset \mathcal{U} \Leftrightarrow v_i \in \mathcal{U} \ominus K\mathcal{E}$$

Set \mathcal{E} known offline – can compute constraints offline!

Ideally \mathcal{E} is the minimum RPI set $\mathcal{F}_\infty = \bigoplus_{j=0}^{\infty} A^j \mathcal{W}$

(c) Convex Optimization Problem

Problem Formulation:

$$\min_{Z,V} l_f(z_N) + \sum_{i=0}^{N-1} l(z_i, v_i)$$

$$\text{s.t. } z_{i+1} = Az_i + Bv_i$$

$$z_i \in \mathcal{X} \ominus \mathcal{E}, \quad u_i \in \mathcal{U} \ominus K\mathcal{E}$$

$$z_N \in \mathcal{X}_f, \quad x(k) \in z_0 \oplus \mathcal{E}$$

$$\text{Control Law : } \mu_{\text{tube}}(x) := K(x - z_0^*(x)) + v_0^*(x)$$

- Optimizing nominal system with tightened state, input constraints
- First tube center z_0 is opt. var.** \rightsquigarrow has to be within \mathcal{E} of x_0
- Cost is w.r.t tube centers, terminal set is w.r.t tightened constraints

Caution: $K(x - z_0^*(x)) + v_0^*(x)$ **NOT LINEAR** in CL

Robust Invariance

Suppose the terminal ingredients ($l_f, \mathcal{X}_f^{\text{ct}}, \pi_f$) are designed such that $\mathcal{X}_f^{\text{ct}} \subset \mathcal{X}$ and for all $z \in \mathcal{X}_f^{\text{ct}}$:

- $\pi_f(z) \in \mathcal{U}$
- $Az + B\pi_f(z) + w \in \mathcal{X}_f^{\text{ct}} \forall w \in \mathcal{W}$
- $l_f(Az + B\pi_f(z)) - l_f(z) \leq -l(z, \pi_f(z))$

Let \mathcal{X}_N be the feasible set and $V^*(x(k))$ be the optimizer of the robust constraint-tightening MPC problem.

Then $Ax(k) + Bv_0^*(x(k)) + w(k) \in \mathcal{X}_N \quad \forall w(k) \in \mathcal{W}$

\rightarrow problem is recursively feasible

Robust Constraint Satisfaction

Tube-MPC Assumptions: almost the same as for nominal MPC

- Stage cost pos def, i.e strictly pos and only 0 at origin
- Terminal set is invariant for the **nominal system** under local control law $\kappa_f(z)$: $Az + B\kappa_f(z) \in \mathcal{X}_f \quad \forall z \in \mathcal{X}_f$
All **tightened state and input constraints** satisfied in \mathcal{X}_f :
 $\mathcal{X}_f \subseteq \mathcal{X} \ominus \mathcal{E}, \quad \kappa_f(z) \in \mathcal{U} \ominus K\mathcal{E} \quad \forall z \in \mathcal{X}_f$
- Terminal cost is cont. Lyapunov function in terminal set \mathcal{X}_f :
 $l_f(Az + B\kappa_f(z)) - l_f(z) \leq -l(z, \kappa_f(z)) \quad \forall z \in \mathcal{X}_f$

Theorem: Robust Invariance of Tube-MPC

Set $\mathcal{Z} := \{x \mid \mathcal{Z} \neq \emptyset\}$ is robust invariant set of system $x(k+1) = Ax(k) + B\mu_{\text{tube}}(x(k)) + w(k)$ subject to constraints $x, u \in \mathcal{X}, \mathcal{U}$

Proof let $\{v_0^* \dots v_{N-1}^*, \{z_0^* \dots z_N^*\}$ be optimal sol'n for $x(k)$ At next point in time, state $x(k+1)$ may have many possible values due to disturbance

By construction, state $x(k+1)$ in the in the set $z_1^* \oplus \mathcal{E} \forall \mathcal{W}$

Therefore the following sequence is feasible for all $x(k+1)$

$$(\{v_1^* \dots v_{N-1}^*, \kappa_f(z_N^*)\}, \underbrace{\{z_1^* \dots z_N^*, Az_N^* + B\kappa_f(z_N^*)\}}_{\text{feas. IC}} \underbrace{\mathcal{X}_f}_{\mathcal{X}_f \rightsquigarrow \text{feas.}})$$

Robust Stability

Robust Stability of Tube-MPC

State $x(k)$ of system $x(k+1) = Ax(k) + B\mu_{\text{tube}}(x(k)) + w(k)$ converges in the limit to the set \mathcal{E}

Proof As in standard MPC we have

$$J^*(z_0^*(x(k))) = l_f(z_N^*) + \sum_{i=0}^{N-1} l(z_i^*, v_i^*)$$

$$\begin{aligned} J^*(z_0^*(x(k+1))) &\leq l_f(z_N^*) + \sum_{i=1}^{N-1} l(z_i^*, v_i^*) \\ &\quad + l(z_0^*, v_0^*) - l(z_0^*, v_0^*) + l_f(z_N^*) - l_f(z_N^*) \\ &= J^*(x(k)) - \underbrace{l(z_0^*, v_0^*)}_{\geq 0} - \underbrace{l_f(z_N^*) + l_f(z_{N+1}) + l(z_N^*, \kappa_f(z_N^*))}_{\leq 0 \text{ (} l_f \text{ is lyap function in } \mathcal{X}_f \text{)}} \end{aligned}$$

This shows $\lim_{k \rightarrow \infty} J(z_0^*(x(k))) = 0$, therefore $\lim_{k \rightarrow \infty} z_0^*(x(k)) = 0$

Caution:

- $x(k)$ does **not tend to 0!** It only stays within robust invar set centered at $z_0^*(x(k))$: $\lim_{k \rightarrow \infty} \text{dist}(x(k), \mathcal{E}) = 0$
- \mathcal{E} must be robust positive invariant for proof (so error remains bounded)

9.3 Tube-MPC Implementation

Offline Design

- Choose stabilizing controller K s.t $\|A + BK\| < 1$
- Compute mRPI set $\mathcal{E} = F_\infty$ for system $x(k+1) = (A+BK)x(k) + w(k), w \in \mathcal{W}$
- Compute tightened constraints $\tilde{\mathcal{X}} := \mathcal{X} \ominus \mathcal{E}, \tilde{\mathcal{U}} := \mathcal{U} \ominus K\mathcal{E}$
- Choose terminal weight function l_f and constraint \mathcal{X}_f satisfying assumptions on tube MPC

LQR Terminal Constraint (typical choice)

- Choose LQR terminal control law $\kappa_f(x) = Kx$, (Q, R same as MPC)
- Find \mathcal{X}_f invar under this controller s.t satisfies constraints

Online Design

- Measure / Estimate state x
- Solve optimization problem:
($V^*(x_0), Z^*(x_0)$) = $\text{argmin}_{V,Z} \{J(Z, V) \mid (Z, V) \in \mathcal{Z}(x_0)\}$
- Set input to $u = K(x - z_0^*(x)) + v_0^*(x)$

Tube-MPC Summary

Benefits

- Less conservative than OL robust MPC (now actively compensating for noise in prediction)
 - Works for unstable systems
 - Optimization problem to solve is 'simple'
- Cons**
- Sub-optimal MPC (optimal extremely difficult)
 - Reduced feasible set when compared to nominal MPC
 - We need to know what \mathcal{W} is (usually not realistic)

9.4 Robust MPC for Uncertain Systems - Summary

Idea compensate for noise in prediction to ensure constraint satisfaction

Benefits

- Feasible set invariant – know exactly when controller will work
- Easier to tune – knobs to tradeoff robustness against performance
- Complex (tubes easy to implement, complex to understand)
- Must know largest noise \mathcal{W}
- Often conservative
- Feas set may be small

9.5 Robust MPC - Extensions

9.5.1 Robust Constraint Tightening MPC

Idea Combine best of Robust OL and Tube-Based MPC

\rightarrow Use propagated error bound to tighten constraints

Error Dynamics:

$$e_{i+1} = (A+BK)e_i + w_i = A_K e_i + w_i, \quad w_i \in \mathcal{W}$$

$$\text{If } e_0 = 0 \text{ then } e_i = \sum_{j=0}^{i-1} A^j w_{i-1-j} \in \mathcal{W} \oplus A_K \mathcal{W} \oplus \dots \oplus A_K^{i-1} \mathcal{W}$$

Problem Setup:

$$\min_{Z,V} l_f(z_N) + \sum_{i=0}^{N-1} l(z_i, v_i)$$

$$\text{subj. to } z_{i+1} = Az_i + Bv_i$$

$$z_i \in \mathcal{X} \ominus (\mathcal{W} \oplus A_K \mathcal{W} \oplus \dots \oplus A_K^{i-1} \mathcal{W})$$

$$u_i \in \mathcal{U} \ominus K(\mathcal{W} \oplus A_K \mathcal{W} \oplus \dots \oplus A_K^{i-1} \mathcal{W})$$

$$z_N \in \mathcal{X}_f \ominus (\mathcal{W} \oplus A_K \mathcal{W} \oplus \dots \oplus A_K^{N-1} \mathcal{W})$$

$$z_0 = x(k)$$

$$\text{Control Law } u(k) = v_0^* + K(x(k) - z_0) = v_0^*$$

Motivation can robustly ensure constraint satisfaction at each time step

Note need terminal set \mathcal{X}_f that is robust invariant under tube controller K

9.5.2 Nominal MPC with Noise

Standard MPC Problem for $x(k+1) = Ax(k) + Bu(k) + w(k)$

$$J^*(x_0) = \min_U l_f(x_N) + \sum_{i=0}^{N-1} l(x_i, u_i)$$

$$\text{s.t. } x_{i+1} = Ax_i + Bu_i, \quad x_i, u_i, x_N \in \mathcal{X}, \mathcal{U}, \mathcal{X}_f$$

Effect on Lyapunov Function

Assume Optimal cost J^* Lipschitz continuous

$$|J^*(Ax + Bu^*(x) + w) - J^*(Ax + Bu^*(x))|$$

$$\leq \gamma \|Ax + Bu^*(x) + w - (Ax + Bu^*(x))\| = \gamma \|w\|$$

Lyapunov Decrease can be bounded as

$$J^*(Ax + Bu^*(x) + w) - J^*(x) = J^*(Ax + Bu^*(x) + w) + J^*(x)$$

$$\leq J^*(Ax + Bu^*(x)) - J^*(x) + \gamma \|w\| \leq -l(x, u^*(x)) + \gamma \|w\|$$

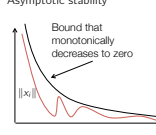
- Amount of decrease grows with $\|x\|$
- Amount of increase upper bounded by $\max\{\|w\| \mid w \in \mathcal{W}\}$

Benefits

- No special knowledge required – "just works" (sometimes)
- Often very effective in practice
- Large feasible set
- Region of attraction may be relative-ly large
- Very difficult to determine region of attraction (set of states where controller works)
- Hard to tune
- Only works for NL systems under continuity assumptions

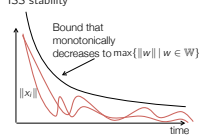
ISS – Input-to-State Stability

Asymptotic stability



System converges to zero

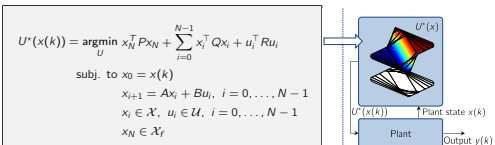
ISS stability



Converges to set around zero, who's size is determined by size of the noise

10 Implementation

10.1 Explicit MPC



Recall: Quadratic Cost State Feedback Solution

MP-QP – Multiparametric Quadratic Program

$$J^*(x(k)) = \min_U \begin{bmatrix} U^T & x(k)^T \end{bmatrix} \begin{bmatrix} H & F^T \\ F & Y^T \end{bmatrix} \begin{bmatrix} U^T & x(k)^T \end{bmatrix}^T$$

$$\text{subj. to } GU \leq w + Ex(k)$$

Solution Properties – $J^*(x(k))$ convex and PW Quad. on polyhedra.

Active Set for $l = 1, \dots, m$

Define active set at $x, A(x)$, and it's complement $NA(x)$ as

$$A(x) := \{j \in l : G_j z^*(x) - S_j x = w_j\} \quad (\text{satisfied with eq.})$$

$$NA(x) := \{j \in l : G_j z^*(x) - S_j x < w_j\} \quad (\text{strict inequality})$$

Critical Region

CR_A is set of parameters x for which set $A \subseteq l$ of constraints i active at the optimum. For given $\bar{x} \in \mathcal{K}^*$ let $(A, NA) := (A(\bar{x}), NA(\bar{x}))$. Then

$$CR_A := \{x \in \mathcal{K}^* : A(x) = A\} \quad (\text{states share active set})$$

Point Location

- Sequential Search** – Computationally linear, very simple, works for all problems
- Search Tree** – Potentially logarithmic, significant offline processing (reasonable for $< 1k$ regions)

Remarks on Explicit MPC

- Linear MPC + Quad / Linear-norm cost \rightsquigarrow Controller PWA func.
- Can pre-compute this function offline
- Online evaluation of PWA function very fast (ns - μ s)
- Can only do this for small systems (3-6 states, small horizon)

10.2 Iterative Optimization Methods

Generic Optimization Problem:

convex if $f : \mathbb{R}^n \rightarrow \mathbb{R}$ and set \mathcal{Q} convex

Analytical sol'n cannot be obtained except simplest cases

$$\text{minimize } f(x)$$

$$\text{subj. to } x \in \mathcal{Q}$$

Iterative Optimization Methods Given initial guess $x^{(0)}$, produce sequence of iterates

$$x^{(i+1)} = \psi(x^{(i)}, f, \mathcal{Q}), \quad i = 0, \dots, m-1$$

such that $|f(x^{(m)}) - f(x^*)| \leq \epsilon$ and $\text{dist}(x^{(m)}, \mathcal{Q}) \leq \delta$ where ϵ and δ are user defined tolerances

10.3 Unconstrained Minimization

Optimality Conditions
Assume $f(\cdot)$ diff'bar at x^* . If f convex, then x^* global min iff $\nabla f(x^*) = 0$

Descent Methods

- $x^{(i+1)} = x^{(i)} + h^{(i)} \Delta x^{(i)}$
- $x^{(i+1)} < f(x^{(i)})$
- Δx : step/search direction
- $h^{(i)}$: step size/length
- $f(x^{(i+1)}) < f(x^{(i)})$ i.e. until termination condition (e.g. $f(x^{(m)}) - f(x^*) \leq \epsilon_1$)
- $\exists h^{(i)} > 0$ s.t. $f(x^{(i+1)}) < f(x^{(i)})$ if $\nabla f(x^{(i)})^T \Delta x^{(i)} < 0$

Descent Direction

- Gradient descent** $x^{(i+1)} = x^{(i)} - h^{(i)} \nabla f(x^{(i)})$
 - Assume ∇f Lipschitz-continuous $|\nabla f(x) - \nabla f(y)| \leq L \|x - y\|$
 - Choose constant step size $h^{(i)} = 1/L$
- Newton Step** $x^{(i+1)} = x^{(i)} + h^{(i)} \Delta x_{nt}$
 - $\Delta x_{nt} = -(\nabla^2 f(x^{(i)}))^{-1} \nabla f(x^{(i)})$
 - Exact Line Search $h^{(i)*} = \text{argmin}_{h>0} f(x^{(i)} + h^{(i)} \Delta x_{nt})$
 - Optimization in 1 var \rightsquigarrow solve by bisection, time consuming
 - Inexact Line search: find $h^{(i)}$ that decreases f by some amount

10.4 Constrained Minimization

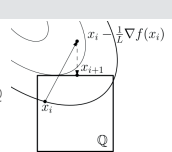
Projected Gradient Methods

Incorporate Constraints in Gradient Step

$$x^{(i+1)} = \pi_{\mathcal{Q}}(x^{(i)} - h^{(i)} \nabla f(x^{(i)}))$$

Projection $\pi_{\mathcal{Q}} = \text{argmin}_x \frac{1}{2} \|x - y\|_2^2$ s.t. $x \in \mathcal{Q}$

- Simple input constraints
- State constraints: hard \rightsquigarrow solve for dual



Interior-Point Methods

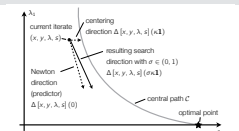
System $\min f(x)$ s.t. $g_i(x) \leq 0, i = 1, \dots, m$ **Assumptions** f, g_i convex, twice cont. diff'bar. $f(x^*)$ is finite and attained, stict feasibility $\exists g(\bar{x}) < 0$, feasible set closed & compact **Idea** Reformulate as unconstrained problem

Primal-Dual Interior-Point Methods

Idea – Iteratively solve relaxed KKT system leave λ^* as variables, linearize and solve resulting system of linear eqns at each iteration

Search Direction $\Delta[x, \nu, \lambda, s](v)$

- $v = 0$ pure Newton direction "predictor"/"affine-scaling"
- $v = \kappa 1$ centering direction, approach \Rightarrow combine via **centering parameter** $\sigma \in (0, 1)$



11 Appendix

11.1 Set Operations

Minkowski Sum: $A \oplus B := \{x + y \mid x \in A, y \in B\}$

Pontryagin Difference: $A \ominus B := \{x \mid x + e \in A \quad \forall e \in B\}$

Caution: $A \ominus B \subseteq A \ominus B \subseteq A$

11.2 Exercises

largest terminal invariant set determine largest terminal invariant set for the system $x_{i+1} = Ax_i + Bu_i$ with constraint $-c \leq u_i \leq c$ under a **stabilizing linear** terminal control law $u(k) = Kx(k)$

- we need $(A+BK)$ stable $\Rightarrow K \in \mathbb{R}_{-}$
- $x \in \mathcal{X} = \{x \mid Kx \in \mathcal{U}\} = \{x \mid \begin{bmatrix} K \\ -K \end{bmatrix} x \leq \begin{bmatrix} c \\ c \end{bmatrix}\}$
- max. control invariant set algorithm:

$$\Omega_0 = \mathcal{X} \rightarrow \text{pre}(\Omega_0) = \left\{x \mid \begin{bmatrix} K \\ -K \end{bmatrix} x_{+1} \leq \begin{bmatrix} c \\ c \end{bmatrix}\right\} = \left\{x \mid \begin{bmatrix} K(A+BK) \\ -K(A+BK) \end{bmatrix} x \leq \begin{bmatrix} c \\ c \end{bmatrix}\right\}$$

$$\Omega_1 = \text{pre}(\Omega_0) \cap \Omega_0 = \left\{x \mid \begin{bmatrix} K \\ -K \end{bmatrix} x \leq \begin{bmatrix} c \\ c \end{bmatrix}\right\} = \Omega_0$$

$\Omega_1 = \Omega_0$ holds always for a stabilizing controller since $|A+BK| \leq 1$ and thus Ω_0 is more restrictive than $\text{pre}(\Omega_0)$

Quadratic Program

- if the problem is convex & Slater's cond. satisfied \Rightarrow KKT are necessary and sufficient for x^* being a global minimizer (convex = only 1 minimum)
- $\min_{x \in \mathbb{R}^n} \frac{1}{2} x^T H x + f^T x$ s.t. $g(x) = Gx \leq 0$ $h(x) = a^T x - b = 0$

- Convexity:** • check if all eigenvalues of the Hessian Matrix H are non-negative • if furthermore $g(x), h(x)$ are linear / affine sets $\Rightarrow \mathcal{X}$ is convex
- Slater's Cond.:** • find a strictly feasible point s.t. Slater's Cond. is satisfied \Rightarrow we have strong duality $p^* = d^*$
- KKT Cond.:** find suitable Lagrange Multipliers λ^*, ν^* s.t. KKT Cond. are satisfied ($\nabla_x L(x^*, \lambda^*, \nu^*) = 0 \Rightarrow \lambda^*, \nu^*$) \Rightarrow then x^* is a global minimizer

1D system: $x(k+1) = x(k) + u(k) + w(k)$
constraints: $\bullet u \in \mathcal{U} = \{u \mid -10 \leq u \leq 10\}$ $\bullet w \in \mathcal{W} = \{w \mid |w| < 1\}$

1) maximum robust invariant set \mathcal{X}_f with control law $u(k) = -x(k)$
CL system: $x(k+1) = w(k)$ with constraint $-x(k) \in \mathcal{U} = \{u \mid -10 \leq u \leq 10\}$

We must be able to recover to the origin after the disturbance, therefore we compute the robust pre-set of $[-10, 10]$: