

Ramsey numbers of books and quasirandomness

Yuval Wigderson

Joint work with David Conlon and Jacob Fox

"Outside of a dog, a book is a man's best friend. Inside of a dog, it is too dark to read."
–Groucho Marx

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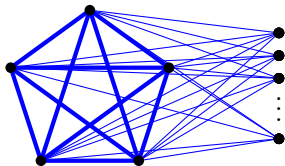
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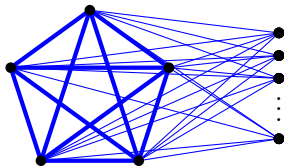
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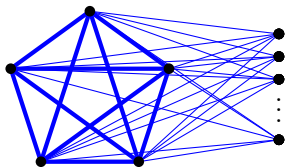
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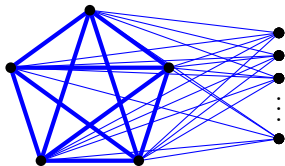


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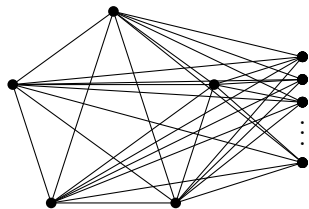
Conlon: $r(K_t) \leq t^{-c \log t / \log \log t} 4^t$.

Book graphs

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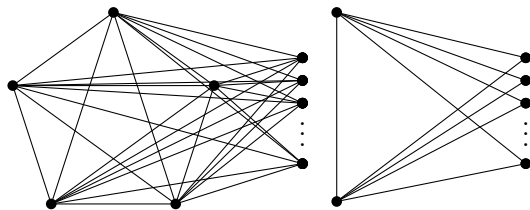
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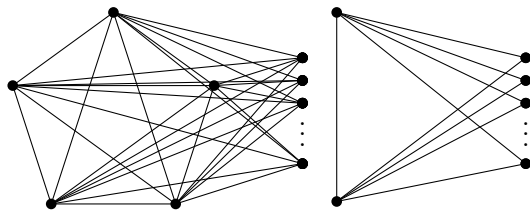
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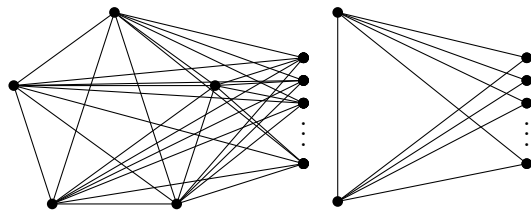
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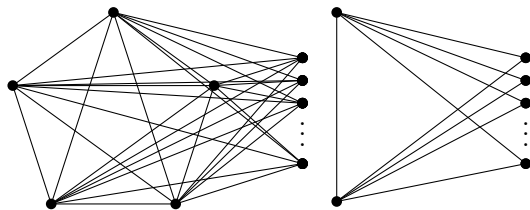


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A coloring of K_N is quasirandom iff at most $o(N^k)$ monochromatic K_k have more than $2^{-k}N + o(N)$ extensions to a monochromatic K_{k+1} .

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For all $\varepsilon > 0$, there exists M such that for every graph G , there is an equitable partition $V(G) = V_1 \sqcup \dots \sqcup V_m$ with $m \leq M$, such that each V_i is ε -regular with itself and with at least $(1 - \varepsilon)m$ other V_j .

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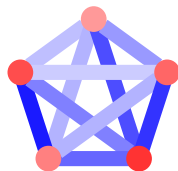
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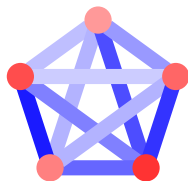
A (k, ε, δ) -**good configuration** is a tuple of disjoint vertex sets C_1, \dots, C_k such that each pair (C_i, C_j) is ε -regular, $d_R(C_i) \geq \delta$, and $d_B(C_i, C_j) \geq \delta$ for $i \neq j$.



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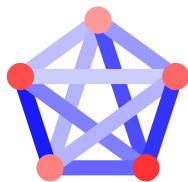
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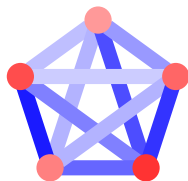


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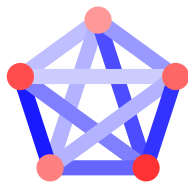
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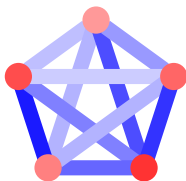
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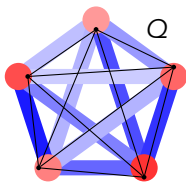
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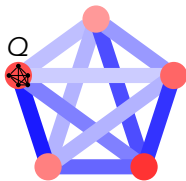
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3. If red, pick $i \in [k]$ uniformly, then a uniform red K_k in C_i .

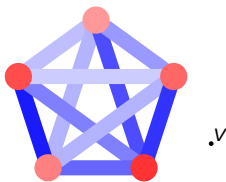
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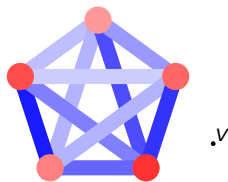
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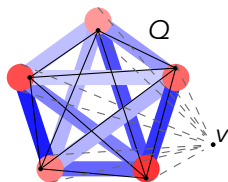
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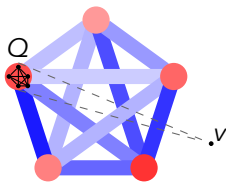
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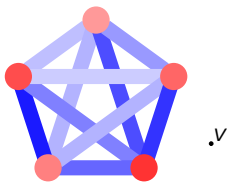
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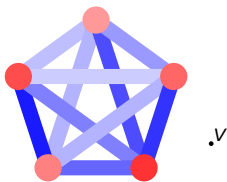
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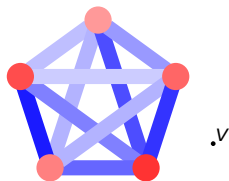
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Key lemma

Lemma (Conlon '18, Conlon-Fox-W)

Let $N = (2^k + \beta)n$ and $\varepsilon \ll \delta \ll \beta$. If a coloring of K_N has a (k, ε, δ) -good configuration, then it contains a monochromatic $B_n^{(k)}$.



Proof sketch:

Let Q be a random monochromatic K_k sampled as above.

For a vertex v , let $x_i(v)$ be the fraction of edges to C_i that are blue.

If C_1, \dots, C_k is a (k, ε, δ) -good configuration,

$$\Pr(Q \cup \{v\} \text{ is monochromatic}) \approx \frac{1}{2} \left(\prod_{i=1}^k x_i(v) + \frac{1}{k} \sum_{i=1}^k (1 - x_i(v))^k \right)$$

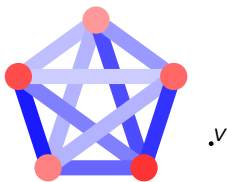
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Adding up over all v shows that Q has $\gtrsim 2^{-k}N$ monochromatic extensions on average \implies monochromatic $B_n^{(k)}$.

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Inductively apply the previous argument to “nibble” out these ε -regular pieces, and conclude quasirandomness. □

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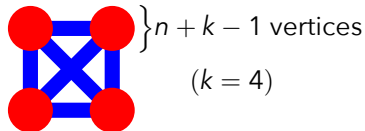
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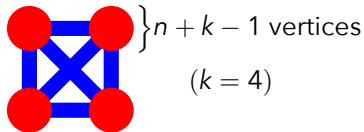
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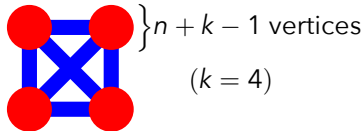
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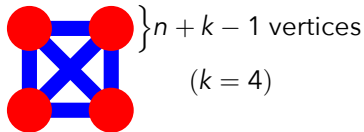
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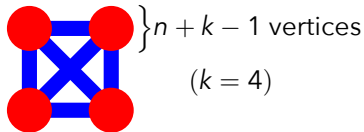
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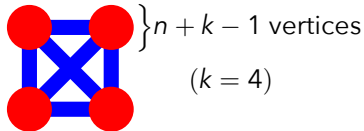
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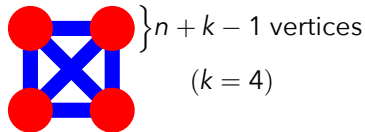
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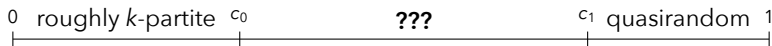


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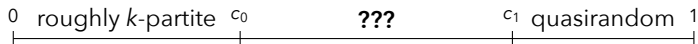
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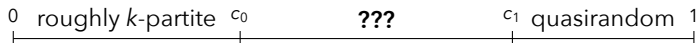
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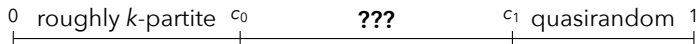
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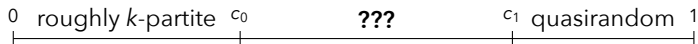
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