

The Ramsey-Theoretic Monster Mash

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ITS Fellows' Seminar

October 31, 2023

Dreadful was the din
Of hissing through the hall, thick swarming now
With complicated monsters head and tail

John Milton, *Paradise Lost* X.521-3

Introduction: behemoths of Ramsey theory

Ghosts of graph Ramsey theory

Sea monsters and Ramsey multiplicity

Shapeshifters and oriented Ramsey numbers

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What is Ramsey theory?



Behemoths

Ghosts

Sea monsters

Shapeshifters

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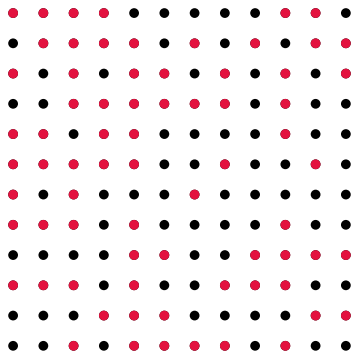
Theorem ("Folklore")

Given N points, if half are colored red, then there are $N/2$ red points.

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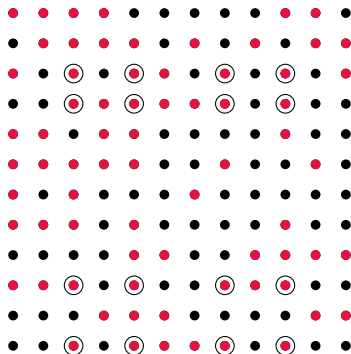
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Given N points, if half are colored red, how many **evenly spaced** red points can we find?



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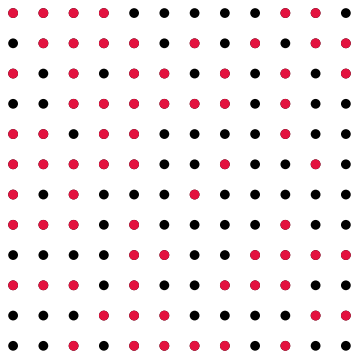
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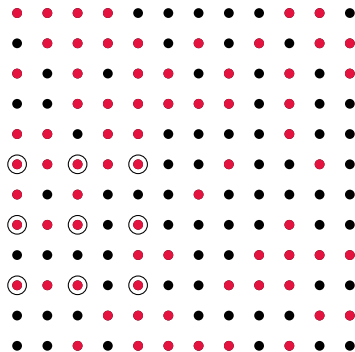
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Such results exist for integers, graphs, posets, Banach spaces...

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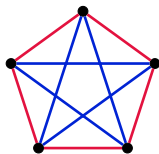
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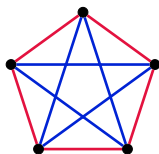
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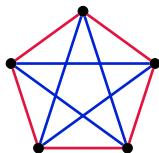


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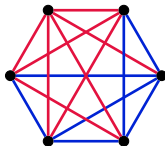
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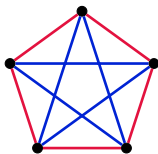
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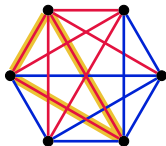
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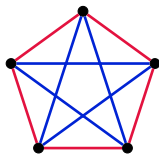
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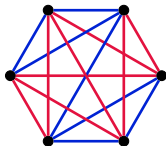
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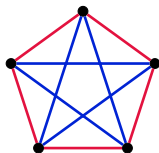
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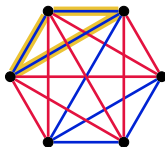
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So there **exists** a coloring of $E(K_N)$ with < 1 monochromatic K_t . \square



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Theorem (Li 2023)

There exists an **explicit** coloring on $N \geq 2^{t^{0.00001}}$ vertices with no monochromatic K_t .



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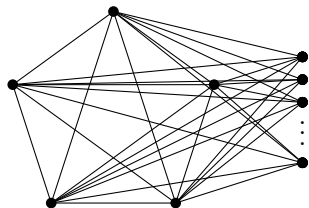


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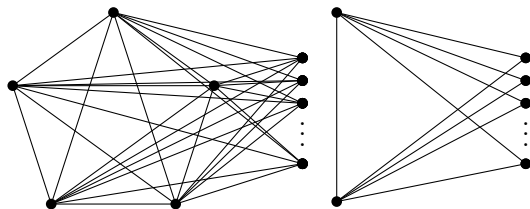


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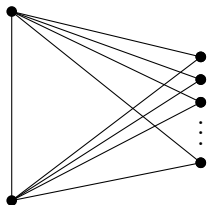
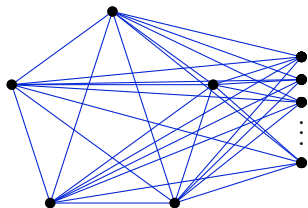


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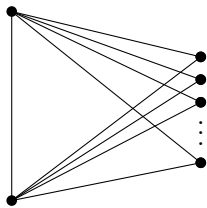
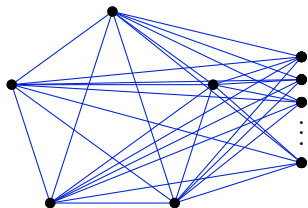


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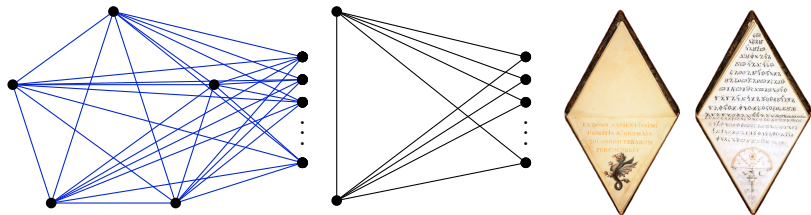


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In the n "page" vertices, it suffices to find a red K_t or a blue K_{t-k} .

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This result is still **far too weak** to improve the bound $r(t) < 4^t$.

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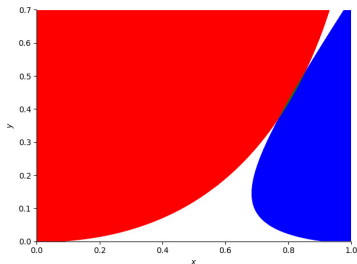
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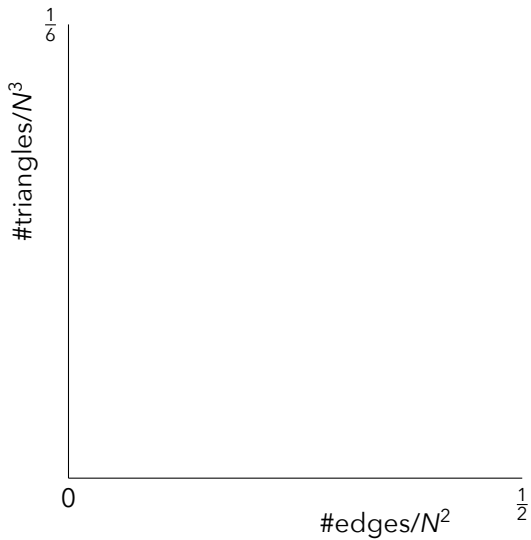
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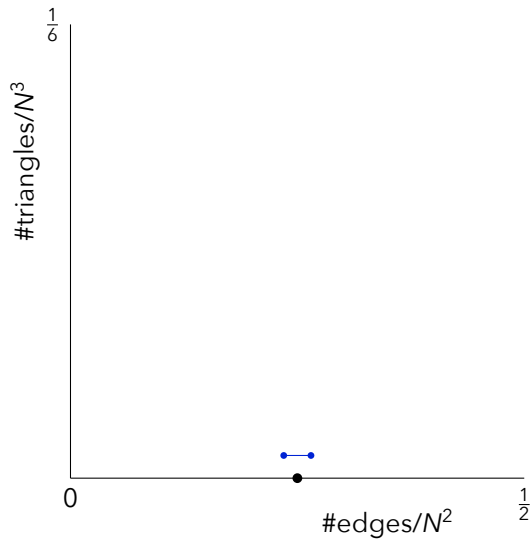
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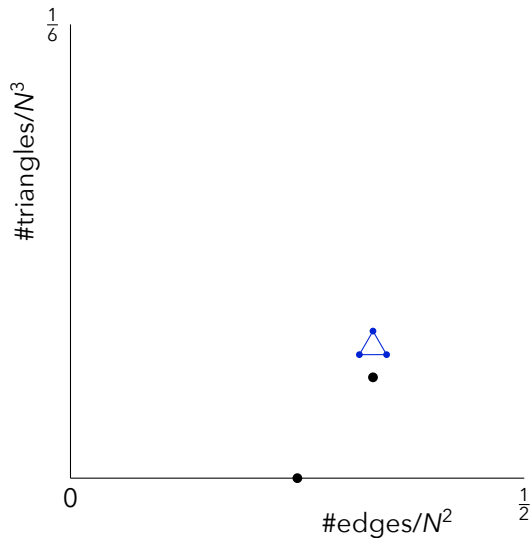
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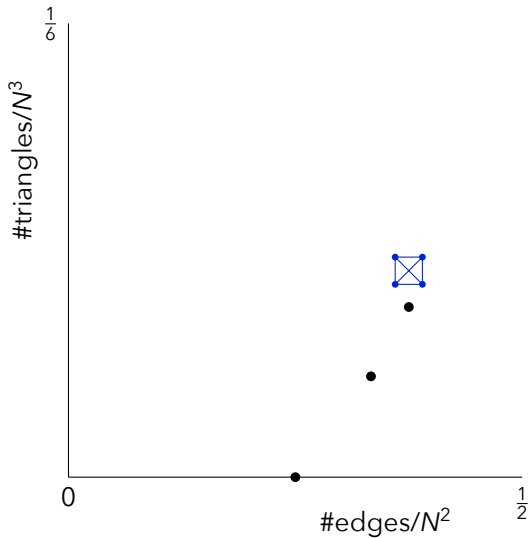
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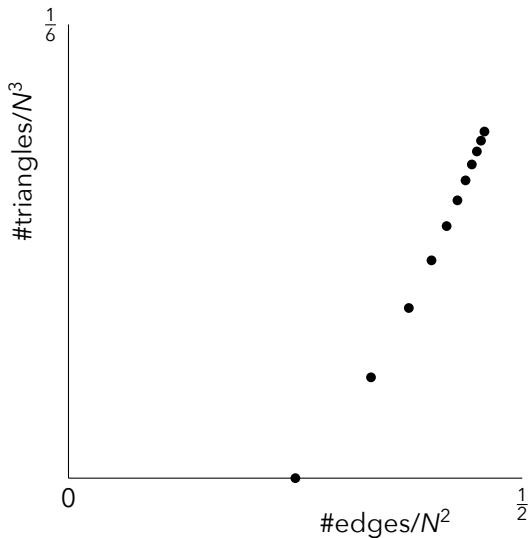
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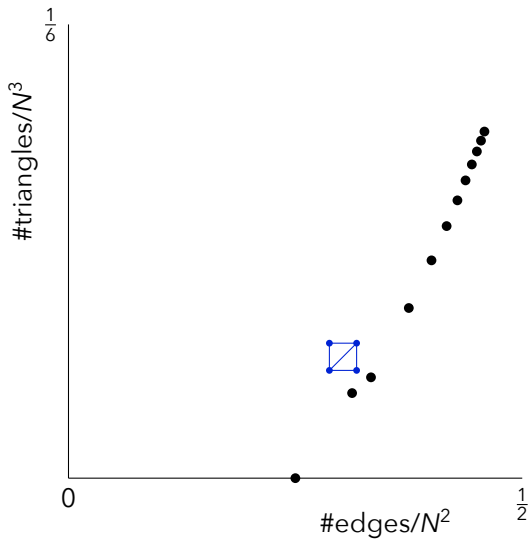
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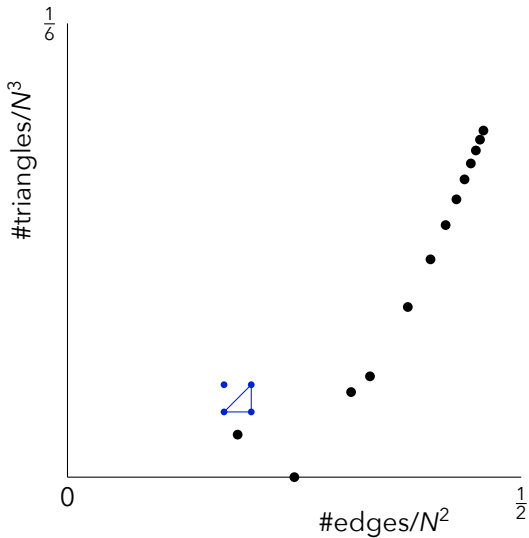
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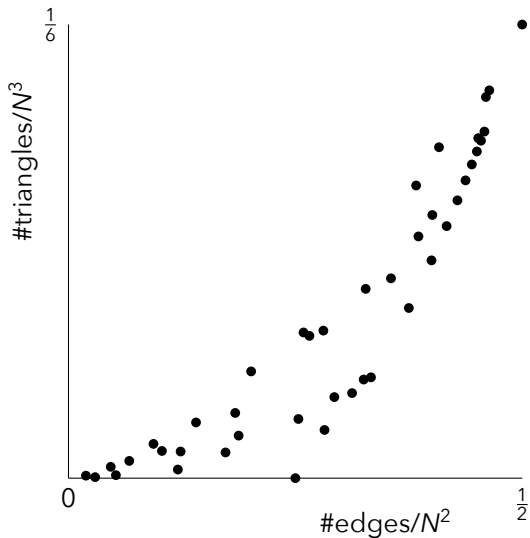
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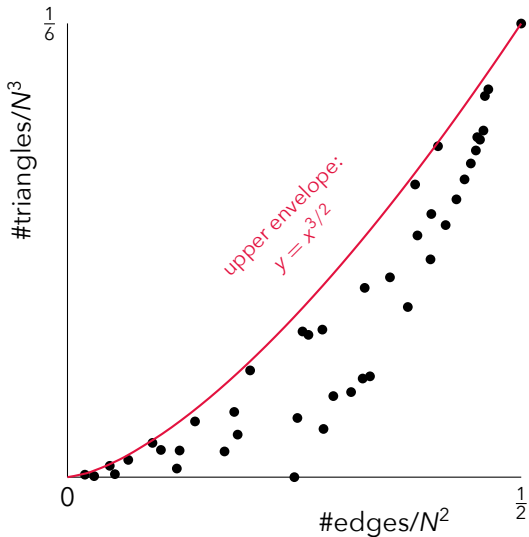
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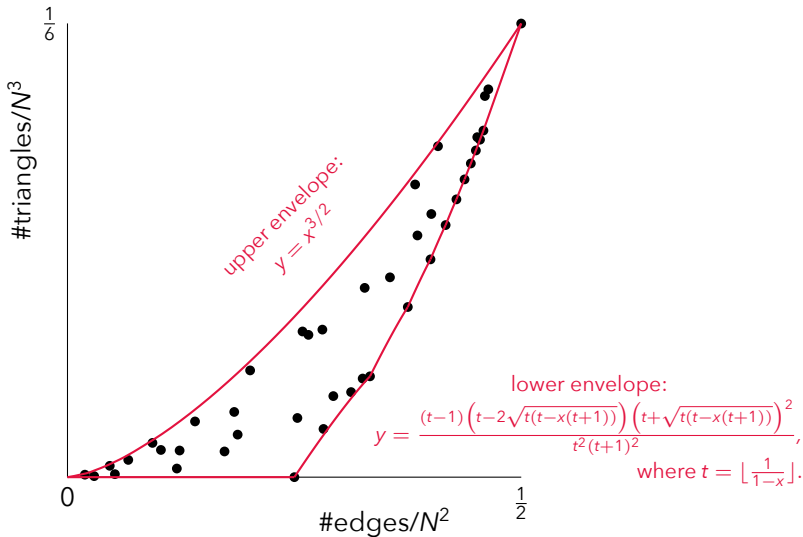
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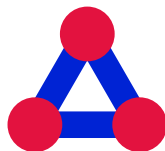


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$k - 1$ parts

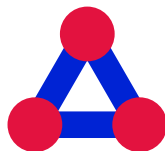


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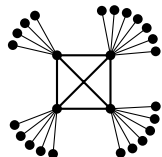
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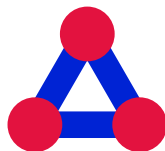


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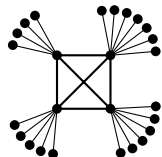
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Open problem: Which coloring minimizes the number of K_4 ?

Introduction: behemoths of Ramsey theory

Ghosts of graph Ramsey theory

Sea monsters and Ramsey multiplicity

Shapeshifters and oriented Ramsey numbers

Ramsey numbers of graphs and digraphs



Ramsey numbers of graphs and digraphs



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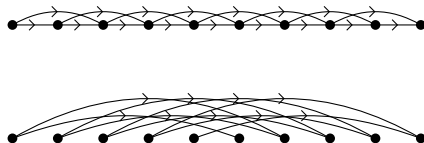


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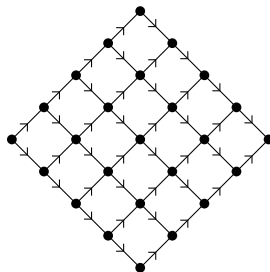
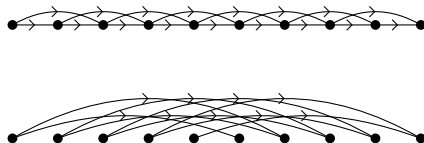


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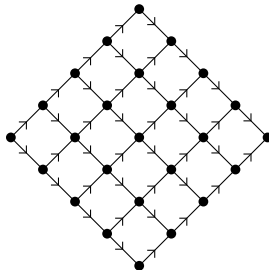
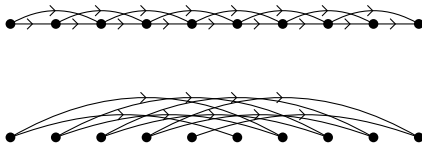


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The digraphs for which $\vec{r}(H)$ is “large” are **shapeshifters**: they have many edges at every length scale, despite having **bounded degree**.

Thank you!