

Recent results on size Ramsey numbers

David Conlon

SIAM Conference on Discrete Mathematics

June 14, 2022

Joint with Jacob Fox and Yuval Wigderson

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Observation (Erdős–Faudree–Rousseau–Schelp)

If H has no isolated vertices, then

$$\frac{r(H)}{2} \leq \hat{r}(H) \leq \binom{r(H)}{2}.$$

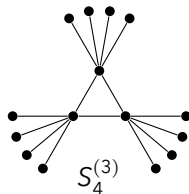
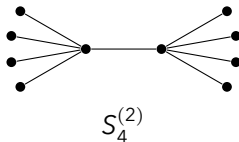
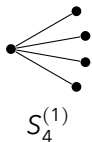
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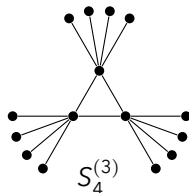
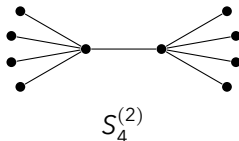
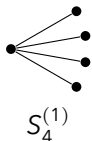
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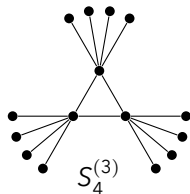
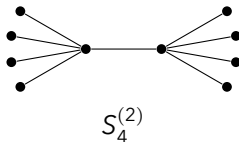
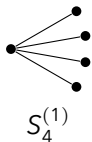


Every n -vertex tree T has $r(T) = \Theta(n)$, so $r(S_n^{(1)})$, $r(S_n^{(2)}) = \Theta(n)$.

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However:

Proposition

$\hat{r}(S_n^{(1)}) = 2n - 1$. On the other hand, $\hat{r}(S_n^{(2)}) = \Theta(n^2)$.

Complete graphs

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Theorem (Erdős, Erdős-Szekeres)

$$2^{t/2} \leq r(K_t) \leq 2^{2t}$$

The four questions

Problem (Erdős-Faudree-Rousseau-Schelp)

Determine the size Ramsey numbers of the following graph families.

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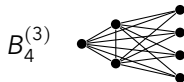
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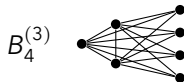
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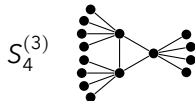
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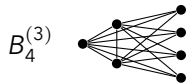
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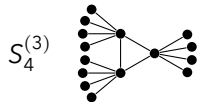
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Any good expander G on $\Theta(n)$ vertices is Ramsey for P_n .

In particular, we may take $G = G(N, p)$ for $N = \Theta(n)$, $p = \Theta(\frac{1}{n})$. \square

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Often, one takes G to be a sparse (pseudo)random graph. To prove that G is Ramsey for H , one uses techniques like sparse regularity.

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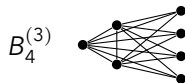
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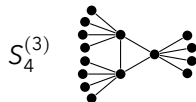
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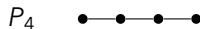
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Pikhurko proved that if $t \gg s$, then this construction (appropriately optimized) is optimal. In particular, $\hat{r}(K_{s,t}) = \Theta(s^2t2^s)$ if $t \gg s$.

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Proof idea: For the lower bound, Erdős-Rousseau use a **uniformly** random coloring. But the upper bound argument is tight if all vertices have **equal** degrees in red and blue.

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So rather than **uniform**, it's better to use a (dyadically iterated) **hypergeometric** random coloring.

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Problem (Erdős-Faudree-Rousseau-Schelp)

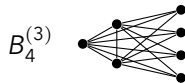
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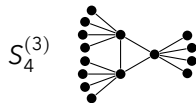
EFRS, ER: $\Omega(st2^s) \leq \hat{r}(K_{s,t}) \leq O(s^2t2^s)$ for $s \leq t$.

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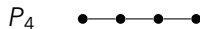
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Lemma (Conlon-Fox-W.)

These are the **only** obstructions to finding monochromatic books.
In particular, $B_N^{(K)}$ is Ramsey for $B_n^{(k)}$ if $N = 2^{k+1} n$, $K = 2kn$.

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The four questions

Problem (Erdős–Faudree–Rousseau–Schelp)

Determine the size Ramsey numbers of the following graph families.

Complete bipartite graphs, $K_{s,t}$

EFRS, ER: $\Omega(st2^s) \leq \hat{r}(K_{s,t}) \leq O(s^2t2^s)$ for $s \leq t$.

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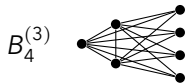
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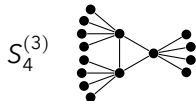
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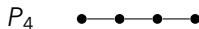
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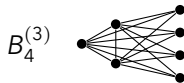
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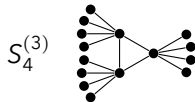
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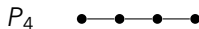
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