## Exercises (recommended)

- 1. (a) Suppose that T is a tree with t+1 vertices, and G is a graph with minimum degree at least t. Prove that G contains a copy of T.

  Hint: Induction on t.
  - (b) Let G be an n-vertex graph with m edges. Prove that there is a subgraph  $G' \subseteq G$  with minimum degree strictly greater than m/n.

    Hint: Repeatedly delete vertices of degree  $\leq m/n$ .
  - (c) Using parts (a) and (b), prove that if T is a tree with t+1 vertices, then

$$ex(n,T) < (t-1)n.$$

(d) Prove that if n is divisible by t, then

$$\operatorname{ex}(n,T) \geqslant \frac{(t-1)n}{2}.$$

? (e) Erdős and Sós conjectured that the lower bound in part (d) is best possible, i.e. that

$$\operatorname{ex}(n,T) = \left| \frac{(t-1)n}{2} \right|$$

for all (t+1)-vertex trees T. Can you prove or disprove this conjecture?

- 2. Let  $K_{1,r}$  denote the star with r leaves. Determine  $ex(n, K_{1,r})$  for all n and r. Is your answer consistent with the Erdős–Sós conjecture from exercise 1? Is it consistent with the Kővári–Sós–Turán theorem we proved in class?
- 3. Recall that we defined

$$m_2(H) = \max_{F \subseteq H} \frac{e(F) - 1}{v(F) - 2},$$

and stated in class that  $ex(n, H) \ge \Omega(n^{2-1/m_2(H)})$  for all bipartite H.

- (a) Compute  $m_2(C_{2\ell})$  for each  $\ell \geq 2$ . What lower bound on  $ex(n, C_{2\ell})$  do you get?
- (b) Compute  $m_2(K_{s,t})$  for all  $t \ge s \ge 2$ . How does the resulting lower bound compare to the others we've discussed?
- (c) Compute  $m_2(T)$  for any tree T. How does the resulting lower bound relate to exercise 1?
- $\star$  (d) Pick your favorite bipartite graph, and compute the lower and upper bounds coming from  $m_2(H)$  and from finding H as a subgraph of  $K_{s,t}$ , respectively. Can you improve either of these bounds?
- 4. Using previous homework problems, prove the following fact. A graph H is a forest if and only if  $ex(n, H) \leq O(n)$ .

 $<sup>\</sup>star$  means that a problem is hard.

<sup>?</sup> means that a problem is open.

 $<sup>\</sup>Leftrightarrow$  means that a problem is on a topic beyond the scope of the course.

## Problems (optional)

- \*1. In this problem, you'll prove that  $ex(n, H) \ge \Omega(n^{2-1/m_2(H)})$ . This problem requires some background in probability, specifically linearity of expectation.
  - (a) Let  $p \in [0, 1]$ , and let G be a random n-vertex graph obtained by making every pair of vertices adjacent with probability p, independently over all choices. Prove that the expected number of edges in G is  $p\binom{n}{2}$ .
  - (b) Prove that for any fixed graph H, the expected number of copies of H in G is at most  $p^{e(H)}n^{v(H)}$ .
  - (c) Suppose that H is 2-balanced, meaning that in the definition of  $m_2(H)$ , the maximizing subgraph F is H itself. Let X denote the random variable defined as the number of edges of G minus the number of copies of H in G. Prove that if  $p = cn^{-1/m_2(H)}$ , for some appropriate constant c > 0, then  $\mathbb{E}[X] \geqslant \Omega(n^{2-1/m_2(H)})$ .
  - (d) Prove that if H is 2-balanced, then  $ex(n, H) \ge \Omega(n^{2-1/m_2(H)})$ .
  - (e) Prove that the same conclusion holds even if H is not 2-balanced.
- $\star 2$ . In this problem you'll prove the Erdős–Sós conjecture in the special case that T is a path. By the *length* of a path, we mean the number of vertices it has.
  - \*\*\*(a) Let G be an n-vertex connected graph with minimum degree  $\delta(G)$ . Prove that G contains a path of length at least  $\min\{n, 2\delta(G) + 1\}$ .

    Hint: Consider a longest path in G, and try to extend it.
    - (b) Let  $P_{t+1}$  denote the path of length t+1. Prove that

$$\operatorname{ex}(n, P_{t+1}) \leqslant \left\lceil \frac{(t-1)n}{2} \right\rceil.$$

*Hint:* Induction on n.

- $\star$  (c) Can you characterize the extremal graphs, i.e. the  $P_{t+1}$ -free graphs with the maximum number of edges?
- 3. Provide an alternative proof of Turán's theorem using induction on r. Let G be a  $K_r$ -free n-vertex graph.
  - (a) Let v be a vertex of maximum degree in G. Let A be the set of neighbors of v, and let  $B = V(G) \setminus A$ .
  - (b) Form a new graph H by deleting all edges inside B, and adding in all missing edges between A and B. Prove that  $e(H) \ge e(G)$ .
  - (c) Apply the inductive hypothesis (Turán's theorem for r-1) to the induced subgraph on A. Conclude that Turán's theorem holds for r.